A Novel Approach for Image Encryption using Dynamic SCAN Pattern

T.Sivakumar and R.Venkatesan

Abstract — In today’s digital era, information is extensively stored and transmitted in the form of digital images, audio, and video files. Security is an important issue for storage and transmission of information due to the growth of electronic data exchange. Digital images are used in fields like telemedicine, military communication, and space research, so it is essential to protect the images from unauthorized disclosure during transmission. In this paper, a novel image encryption method is proposed using scan pattern, circular shift, and transposition methods. The pixels of the original image are permuted by the scan pattern generated by the notion of $K^{th}$ smallest and stack dynamically. The circular shift and transposition are done by Shuffling Key (SK) generated from the original image. Value transformation is performed to confuse the relationship between the original and encrypted images using bitwise XOR operation. The proposed method satisfies the performance analysis tests such as histogram, correlation, large key space, and acceptable encryption speed. Also, the proposed method is resistant to statistical, brute force, and differential attacks.

Index Terms — Information Security, Image Encryption, Dynamic SCAN Patterns, Pixel Permutation

I. INTRODUCTION

SECURITY and ethical issues in information systems are important concerns for most organizations and today’s Internet communication. New media technologies such as social networking and video sharing are used by individuals and organizations to share information [1]. The phenomenal growth of Internet in the recent years has brought privacy concerns to the fore front. Data privacy, data management, data security, and protection of systems from hackers are very important for any organization [2]-[3].

Cryptography is an art of achieving security by encrypting messages to make them non-readable at the sender side and decrypting the messages at the receiver end to obtain the original information. The conventional cryptosystems such as DES – Data Encryption Standard, AES - Advanced Encryption Standard, and IDEA - International Data Encryption Algorithm have been designed to protect textual data and may not be suitable for multimedia data. The use of conventional cryptosystems to encrypt images directly is not adequate for two reasons [4]: (a) the image size is much larger than the text and it needs more time to encrypt the images and (b) the decrypted text must be equal to the original text, but this is not required for images, that is, small distortion in the decrypted image is acceptable by human perception system.

A scanning of a two dimensional array is an order in which each element of the array is accessed exactly once. This is a formal language-based two dimensional spatial accessing methodology which could generate a large number of scanning paths [5]. An image could be encrypted in three ways namely, (a) bit permutation, (b) pixel permutation, and (c) block transformation. In this paper, a novel image encryption method using pixel permutation and transposition is proposed.

The rest of the paper is organized as follows. Section two reviews existing image cryptosystems and section three presents the proposed image encryption method with an example. Section four gives the experimental results and analysis. Section five and six deals with the real time implementation choices and conclusion of the paper.

II. RELATED WORK

Han Shuihua et al. [4] designed an asymmetric image encryption scheme based on matrix transformation. This scheme is based on matrix transformation, which is easily implemented and highly efficient to quickly encrypt and decrypt images. Maniccam et al. [5] revealed an image and video encryption method based on SCAN methodology, which is a formal language-based two-dimensional spatial accessing technique. The basic idea of this image encryption method is to rearrange the pixels of the image and change the pixel values. The pixel rearrangement is done by scan keys and the pixel values are changed by a simple substitution mechanism.

The intelligible information present in an image is due to the correlations among the bits, pixels and blocks in a given arrangement. This perceivable information could be reduced by decreasing the correlation among the bits, pixels and blocks using certain combinational permutation techniques. The permutation of pixels and blocks are good at producing higher level security compared to bit permutation [6]. In Hill cipher, if the key matrix is not invertible, then encrypted text could not be decrypted. Bibhudendra et al. [7] presented an Advanced Hill encryption technique, which uses an involutory matrix, to encrypt images. The complexity could be reduced by avoiding the process of finding inverse of the matrix during decryption.

Chin Chen et al. [8] described an image encryption algorithm based on Vector Quantization to encrypt many images with only one codebook. Guodong Ye [9] proposed an image encryption using double logistic maps, in which the digital image is confused from row and column respectively. Confusion and diffusion effects are carried out solely by the substitution and diffusion stages. Chen’s system is employed to diffuse the gray value distribution. Haojiang Gao et al. [10] presented a Nonlinear Chaotic Algorithm (NCA) using power and tangent functions instead...
of linear function. This encryption algorithm is a one-time-one-password system.

Kanso et al. [11] suggested an algorithm based on three-dimensional chaotic maps. There are three rules to determine the shuffling, mixing and scrambling processes of the pixel values in the plain-image. The image pixels are shuffled according to a search rule based on the 3D chaotic map. Khaled Loukhaoukha et al. [12] brought an image encryption algorithm based on Rubik’s cube principle. The original image is scrambled using the principle of Rubik’s cube and then to confuse the relationship between original and encrypted images where XOR operation is applied to rows and columns of the scrambled image. Liu Hongjun et al. [13] designed a stream cipher based on one-time keys and robust chaotic maps. The algorithm uses piecewise linear chaotic map as the generator of a pseudo-random key sequence.

Mohammad Ali et al. [14] introduced an image encryption algorithm based on block-based transformation and the conventional Blowfish algorithm. The algorithm resulted in lowest correlation and the highest entropy when the transformation algorithm and Blowfish are combined together. It is also found that an inverse relationship exists between the number of blocks and correlation, and a direct relationship between the number of blocks and entropy. Sanfu Wang et al. [15] presented an image scrambling method based on a folding transform with folding matrix which is orthogonal thereby enabling to fold images either up-down or left-right. The algorithm has effective hiding ability for image information with matrix addition and subtraction. Sathishkumar et al. [16] developed an image encryption technique using multiple chaotic based circular mapping. The image is encrypted using logistic map, sub key and its transformation leads to diffusion process. Image encryption based on random pixel permutation and random shuffling strategies are used for permuting pixel coordinates [17]-[18].

Tzung Her Chen et al. [19] proposed a Random Grids-based Visual Secret Sharing scheme with the capability of encrypting multiple secret images at once into only two circular cipher-grids. To decrypt all secrets, decoders stack the two circular cipher-grids to disclose the first secret and then gradually rotate one circular cipher-grid at a fixed degree to disclose the second image. Wenping Guo [20] presented an image scrambling encryption algorithm based on chaotic sequence. The algorithm utilizes the good features of chaotic sequence related to cryptographic properties, such as pseudo-random and sensitivity to initial conditions. Xiaomin Wang et al. [21] presented an image scrambling method using Poker shuffle, which is controlled dynamically by chaotic system. The pixel positions are scrambled by shuffle sequences, which are generated by a chaotic-controlled Poker shuffle process.

Thus, image encryption based on chaotic sequence, scan pattern, permutation, pseudorandom numbers, and block transformation are applied by researchers in the literature. Nevertheless, it is imperative to make use of the concepts of $k^n$ smallest and stack to encrypt images. Hence, in this paper, an attempt is made to encrypt digital images using the above said concepts.

III. THE PROPOSED SYSTEM

The input image($X$) of size $mn$ is taken and the pixel positions are rearranged as follows. The 1st smallest pixel element of each row is scanned and pushed onto the stack. After scanning the 1st smallest pixel in the last row, the content of the stack is popped to form the first row of cipher matrix. Then the 2nd smallest pixel of each row is scanned and pushed onto the stack. After scanning the 2nd smallest pixel of the last row, the content of the stack is popped to form the second row of cipher matrix. This process is continued until scanning the $n$th smallest pixel in all the rows. While scanning the image, the column order in which the first row is accessed is stored as shuffling key. This shuffling key is used for circular shift (for both row and column), permutation (for both row and column), and bitwise XOR operation. The encrypted image is transmitted along with the SCAN key and shuffling key. The encryption process of the proposed method is shown in Figure 1.

A. Pseudocode of Encryption Function

The encryption function performs tasks such as (a) encrypt the image (b) generate SCAN key for permutation, and (c) generate Shuffling Key (SK) to perform circular shift, transposition and XOR operations. The decryption is just the inverse of the encryption process.

Input: Original Image of size rowsize x colsiz
Output: Cipher image of size rowsize x colsiz, Scan key and the Shuffling key

1. Let original_image[rowsize, colsiz] be the original image for encryption, where rowsize, and colsiz represent the number of rows and columns, respectively.
2. Initialize i, j, k, top equal to 1;

```python
Pseudocode of Encryption Function

1. Initialize i, j, k, top equal to 1;
2. While i < rowsize and j < colsiz:
   a. Get the smallest pixel Y = min(plaintext[i, j])
   b. Push Y onto the Stack
   c. Increment i, j, k
3. While top > 0:
   a. Pop Y from the Stack
   b. Place Y at position (i, j)
   c. Increment i, j, k
4. Return Y
```

Fig. 1. Typical Encryption Process of Proposed System
3. Repeat for \( i::1 \) to rowsize,
   (a) Repeat for \( j::1 \) to colsize,
       i. \( \text{stack}[\text{top}] \leftarrow \text{Min}(\text{original_image}[i,:]) \), such that minimum pixel value for the \( i^{th} \) row can be pushed on to stack.
       ii. Increment top;
   (b) Repeat until the stack is empty
       i. \( \text{i_image}[:, \text{k}] \leftarrow \text{stack}[\text{top}] \), such that an intermediate cipher image is created by arranging the minimum pixel values from each row.
       ii. \( \text{scan_key}[i, \text{k}] \leftarrow \text{index}(\text{stack}[\text{top}]) \), such that the index of each minimum pixel value for the \( i^{th} \) row is popped and assigned as scan key for permutation.
       iii. Decrement top and increment \( \text{k} \);
   (c) Increment \( \text{i} \);
4. Assign \( \text{shuffling_key} \leftarrow \text{scan_key}[1, \text{i}] \);
5. Repeat for \( i::1 \) to colsize
   \( \text{R_Shift}[i, :] \leftarrow \text{row_circularshift} \) (original_image, \( \text{shuffling_key}[i] \)), such that each row is right shifted circularly.
6. Repeat for \( i::1 \) to rowsize
   \( \text{Col_Shift}[:, \text{i}] \leftarrow \text{col_circularshift}(\text{R_Shift}, \text{shuffling_key}[i]) \), such that each column is shifted down circularly.
7. Repeat for \( i::1 \) to colsize // Row transposition
   \( \text{Row_Trans}[i, :] \leftarrow \text{row_trans}(\text{shuffling_key}[i]) \),
8. Repeat for \( i::1 \) to rowsize // Column transposition
   \( \text{Col_Trans}[:, i] \leftarrow \text{col_trans}(\text{shuffling_key}[i]) \)
9. Repeat for \( i::1 \) to colsize
   \( \text{Row_XOR}[:, i] \leftarrow \text{bitxor}(\text{Col_Trans}[:, i], \text{shuffling_key}[i]) \), such that each row elements are XORed with the shuffling_key.
10. Repeat for \( i::1 \) to rowsize
    \( \text{Cipher_Matrix}[:, i] \leftarrow \text{bitxor}(\text{Row_XOR}[:, i], \text{shuffling_key}[i]) \), such that each column elements are XORed with the shuffling_key.
11. Return the Cipher_Matrix.

B. Illustration

In this section, the working model of the proposed system is illustrated using an image of size 5X5 pixels.

Step 1: Input the plain image (\( x \)).
Let us consider the following gray scale image shown in Figure 2(a). The intensity value of each pixel is shown in Figure 2(b).

![Original Image](a)

```
240 32 65 192 128
45 86 231 93 141
255 100 55 200 137
115 250 89 46 10
56 125 139 25 225
```

![Gray Value of Original Image](b)

Step 2: Read the 1st smallest pixel value from each row and push into the stack. Once all the rows are scanned for the smallest element, the elements are popped from stack to form the first row of the cipher matrix. The output of this process is shown in Figures 3(a) and (b).

The first smallest value in row one is at index two and it is stored as the first shuffling key. Figures 3(c) and (d) show the content of Shuffling Key (SK) array and SCAN key array after the first iteration.

![Content of Stack](a)

```
25 10 55 45 32
```

![First Row of the Cipher Matrix](b)

```
2
```

![Shuffling Key](c)

```
1, 2, 2, 1, 3, 3, 4, 5, 5, 4
```

![SCAN Key](d)

Fig. 3. (a) – (d) Output after First Iteration

The above step is repeated for 2nd, 3rd, 4th and 5th smallest in each row. The content of cipher matrix, the corresponding image and the value of the shuffling key array are shown in Figure 4(a), (b) and (c), respectively.

![Matrix after SCAN](a)

```
2 3 5 4 1
```

![Image after SCAN](b)

```
25 10 55 45 32
56 48 100 86 65
125 89 137 93 128
139 115 200 141 192
225 250 255 231 240
```

![Shuffling Key](c)

Fig. 4. (a) – (c) Output after SCAN

Figure 5(a) shows the content of SCAN key array after scanning all elements of the original matrix and Figure 5(b) is the reduced SCAN key matrix after removing the column numbers. Since, the first column of the reduced scan key array is same as the shuffling key it is enough to send the reduced 5x5 matrix as SCAN key and Shuffling Key (SK) for decryption.

![SCAN Key](a)

```
1, 2, 2, 1, 3, 3, 4, 5, 5, 4
```

![Reduced SCAN Key](b)

Fig. 5. (a) SCAN Key

Step 3: Apply circular shift (for both row and column wise) using the shuffling key generated in step (2). The row and column wise circular shift is performed as follows: The first row is shifted right by two pixel positions using the shuffling key 2, with the first element of the row circularly following the last. Likewise, the second row is shifted right by three pixel positions using the shuffling key 3 and this is
done for the 3rd, 4th, and 5th rows using the shuffling keys 5, 4, and 1 respectively. The same process is repeated for columns using the same shuffling key. That is, the first column is shifted two pixel positions beneath, with the top element of the column circularly following the last. The final result of step (3) is shown in Figures 6(a) and (b).

Step 4: Perform row and column permutation by using the shuffling key generated in step (2). The row and column wise permutations are performed by rearranging the columns first and then rows according to the order 2, 3, 5, 4 and 1. Figures 7(a) and (b) show the result after permuting rows and columns.

Step 5: Perform row wise bitwise XOR operation. Figures 8(a) and (b) show the output after performing bitwise XOR operation for all rows and columns using the shuffling key. Result shows that the encrypted matrix is completely different from the original pixel matrix shown in Figure 2(b).

The result shows that the encrypted image is different from the plain image and the correlation between the plain and encrypted images is close to zero. The correlation value between the original and encrypted images is 0.0029 for SCSP and -0.0139 for SCSPX.

A. Analysis of Histogram

Histogram is a graphical representation to show the visual impression of the distribution of data. An image histogram is a graphical representation of the pixel distribution in an image. It plots the number of pixels for each pixel value [22].

Figure 10 shows the histogram of the Lena and Map images before and after encryption. The histogram of the encrypted image shows that the gray scale values are uniformly distributed in the encrypted image. Hence the proposed method is safe from histogram analysis attacks and satisfies the diffusion property.
B. Analysis of Correlation

Correlation is a statistical analysis technique that could show whether and how strongly pairs of variables are related [22]. Correlation determines the relationship between the original image and cipher image. It is a measure that computes degree of similarity between the images and it is a useful measure to judge the encryption quality of cryptosystems. The image cryptosystem is said to be good, if encryption algorithm hides all attributes of the plaintext image, encrypted image is totally random and highly uncorrelated. This is calculated by the equation (1).

\[ \gamma_{xy} = \frac{\text{Cov}(x, y)}{\sqrt{D(x)D(y)}} \]  

\[ D(x) = \frac{1}{N} \sum_{i=1}^{N} (x_i - E(x))^2 \]

\[ E(x) = \frac{1}{N} \sum_{i=1}^{N} x_i \]

\[ \text{Cov}(x, y) = \frac{1}{N} \sum_{i=1}^{N} (x_i - E(x))(y_i - E(y)) \]

Where, \( x_i \) and \( y_i \) are the gray scale values of corresponding pixels from the original and encrypted images, \( N \) is the number of pixel pairs \( (x_i, y_i) \), \( E(x) \) and \( D(x) \) are the mean and standard deviation of the gray level values.

Figure 11 shows the correlation value between (a) original Vs. encrypted images using SCAN pattern; denoted as SCAN, (b) original Vs. encrypted images using SCAN and circular shift; indicate as SCS, (c) original Vs. encrypted images using SCAN, circular shift, and permutation; represented as SCSP, and (d) original Vs. encrypted images using SCSP and bitwise XOR operation; denoted as SCSPX.

The correlation value shows that the statistical relationship between the original and encrypted images varies between 0.4 to -0.4. The correlation value is close to zero after the operations SCAN, circular shift, and permutation. Figure 12 (a) – (f) shows correlation distributions of the horizontal, vertical, and diagonal adjacent pixels of the original baboon image and the corresponding encrypted image.
Fig. 12. Adjacent Pixels Correlation Distribution
Table II shows the corresponding correlation value and the correlation value is reduced considerably for the encrypted image. Hence, the proposed system satisfies confusion property significantly.

<table>
<thead>
<tr>
<th></th>
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<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>Lena</td>
<td>0.985</td>
<td>0.973</td>
<td>0.962</td>
<td>0.342</td>
<td>0.352</td>
<td>0.298</td>
</tr>
<tr>
<td>Baboon</td>
<td>0.975</td>
<td>0.983</td>
<td>0.972</td>
<td>0.293</td>
<td>0.395</td>
<td>0.207</td>
</tr>
<tr>
<td>Deer</td>
<td>0.969</td>
<td>0.958</td>
<td>0.948</td>
<td>0.249</td>
<td>0.386</td>
<td>0.174</td>
</tr>
<tr>
<td>Pepper</td>
<td>0.980</td>
<td>0.979</td>
<td>0.971</td>
<td>0.296</td>
<td>0.324</td>
<td>0.266</td>
</tr>
</tbody>
</table>

The results were taken after a single iteration of the proposed scheme. The correlation value after performing the proposed scheme on Baboon, Cameraman, and Planet images for the number of iterations varying in terms of five is shown in Figure 13.

The correlation value is rapidly decreased for five iterations and for subsequent iteration, there is no drastic change. Thus, it is enough to repeat the SCSP process for five to ten iterations to reduce the processing overhead.

C. Cut Test Analysis

The prohibited users may cut or destroy the information condition so that they could not be decrypted. The cut test analysis is performed using Lena image. Figures 14(a) - (c) show the encrypted Lena images after removing 25x25, 50x50, and 75x25 pixel regions, respectively. Figures 14(d) - (f) show the corresponding decrypted images.

The decrypted images are negligibly distorted and could be recognized by human perception system as a Lena image. The proposed method is harmless from cropping attack and loss of data during transmission. It is also seen that the removed region is spread all over the image, and hence the proposed scheme has good dispersion rate.

D. Sensitivity of Shuffling Key

For successful decryption, the Shuffling Key (SK) is an important component. The proposed method is tested for the effect of small changes in the SK and the effect of change in SK value on Baboon image is shown in Figure 15.

The result shows that the decrypted image is completely different from the original Baboon image, and hence the proposed method is very much sensitive with respect to the shuffling key value. i.e., a small change in SK will degrade and affect the quality of decrypted image.

E. Key Space Analysis

The elapsed time taken for encryption, including SCAN key generation, is approximately 12 sec and decryption is 0.62 sec. If the scan key in unknown, then there are n! combinations to place the pixels of the first row, where n is the number of columns. An image of size m*n pixels will have m^n! combinations to place all pixel elements of the image. Hence, the key space of an image of size 256x256 pixels is $10^{509}$. If one decryption takes 0.5 sec, then an attacker need, approximately, $3.49e^{501}$ years for exhaustive key search attack. The time complexity of encryption() and decryption() functions are O(mxn).

V. REAL TIME IMPLEMENTATION STRATEGIES

The SCAN and Shuffling Keys are generated from the plain image, and the proposed cryptosystem is a one-time-one-password scheme. A key image could be used to generate the scan pattern and shuffling key to encrypt and decrypt other images in real time implementation. The key image is shared securely before communication and this will eliminate overhead of transmitting the scan and shuffling keys for each image. If any suspicion or security breach is found in the network, then the key image could be altered. Figures 16(a) - (e) show the results of step by step encryption of Lena image using Baboon as a key image. Figure 16(f) shows the histogram of the encrypted Lena image and it is very much flat.
Figure 16(g) - (i) shows the correlation distribution between adjacent pixels of encrypted Lena image. The result shows that the proposed method satisfies the diffusion property and adjacent pixels in encrypted image are weakly correlated.

A. Results of Dispersion Test

Figures 17(a) - (f) illustrate the results of dispersion test using a white image with 15x15 black region. The results show that the 15x15 region is distributed across the entire image, and hence the proposed scheme has acceptable pixel permutation.
The proposed method could be implemented using any one of the following cases in real time applications. Case (1): applying SCAN pattern and XOR operation, case (2): applying SCAN pattern, row and column transposition, and XOR operation, and case (3): applying SCAN pattern, row and column circular shift, and XOR operation.

B. Results of Case 1

The result of case (1) using Baboon as key image and Lena as plain image is shown in Figure 18. Result shows that the encrypted image is different from the original image and thus, the proposed scan pattern provides acceptable pixel permutation. The expected time taken for encryption and decryption is 0.0382 and 0.0330 seconds, excluding SCAN key generation, respectively. The correlation between the original and encrypted images is 0.0042.

C. Results of Case 2

The result of case (2) using Baboon as key image to encrypt Lena image is shown in Figure 19. Results show that the encrypted image in each step is entirely different from the original Lena image and hence, the combination of SCAN, transposition, and XOR offer reasonable hiding. The encryption and decryption time is 0.0140 and 0.0438 seconds, except SCAN key generation, respectively. The correlation between the original and encrypted images is -0.0013.

D. Results of Case 3

The results of case (3) using Baboon as key to encrypt Lena image is shown in Figure 20. The combination of proposed SCAN and circular shift could provide acceptable encryption result. The estimated time taken for encryption and decryption is 0.4532 and 0.5365 seconds, excluding SCAN key generation, respectively. The correlation between the original and encrypted images is -0.0089.
E. Adjacent Pixel Correlation

The horizontal, vertical, and diagonal correlation values of the encrypted images are given in Table III. The correlation value is reduced considerably after encryption and therefore, the proposed system satisfies confusion property significantly.

<table>
<thead>
<tr>
<th>Case</th>
<th>Directions</th>
<th>Lena</th>
<th>Baboon</th>
<th>Cameraman</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Horizontal</td>
<td>0.1610</td>
<td>0.1316</td>
<td>0.1394</td>
</tr>
<tr>
<td></td>
<td>Vertical</td>
<td>0.1999</td>
<td>0.1618</td>
<td>0.1552</td>
</tr>
<tr>
<td></td>
<td>Diagonal</td>
<td>0.0548</td>
<td>0.0211</td>
<td>0.0486</td>
</tr>
<tr>
<td>2</td>
<td>Horizontal</td>
<td>0.1915</td>
<td>0.1522</td>
<td>0.1923</td>
</tr>
<tr>
<td></td>
<td>Vertical</td>
<td>0.1739</td>
<td>0.1445</td>
<td>0.2718</td>
</tr>
<tr>
<td></td>
<td>Diagonal</td>
<td>0.0452</td>
<td>0.1100</td>
<td>0.1159</td>
</tr>
<tr>
<td>3</td>
<td>Horizontal</td>
<td>0.1686</td>
<td>0.1203</td>
<td>0.2449</td>
</tr>
<tr>
<td></td>
<td>Vertical</td>
<td>0.1645</td>
<td>0.1627</td>
<td>0.2653</td>
</tr>
<tr>
<td></td>
<td>Diagonal</td>
<td>0.0051</td>
<td>0.0269</td>
<td>0.1258</td>
</tr>
</tbody>
</table>

The horizontal, vertical, and diagonal correlations between adjacent pixels for a few existing image encryption algorithms are given in Table IV. The result shows that the proposed method correlation value is approximately close to the correlation value of existing algorithms mentioned in Table IV.

<table>
<thead>
<tr>
<th>Encryption Methods</th>
<th>Horizontal</th>
<th>Vertical</th>
<th>Diagonal</th>
</tr>
</thead>
<tbody>
<tr>
<td>Ref [10]</td>
<td>-0.0159</td>
<td>-0.0654</td>
<td>-0.0323</td>
</tr>
<tr>
<td>Ref [12]</td>
<td>0.0068</td>
<td>0.0091</td>
<td>0.0063</td>
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<td>Ref [14] Case 1</td>
<td>0.034</td>
<td>0.038</td>
<td>0.013</td>
</tr>
<tr>
<td>Ref [14] Case 2</td>
<td>0.0245</td>
<td>0.0067</td>
<td>0</td>
</tr>
<tr>
<td>Ref [14] Case 3</td>
<td>0.0129</td>
<td>0.0034</td>
<td>0.0014</td>
</tr>
<tr>
<td>Ref [23]</td>
<td>0.0263</td>
<td>0.0163</td>
<td>0.0114</td>
</tr>
<tr>
<td>Ref [24]</td>
<td>0.01776</td>
<td>0.04912</td>
<td>0.00348</td>
</tr>
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</table>

F. Randomness Test

To resist the differential attacks, the cipher image should be sensitive to small changes in the plain image, and the cipher image must have random pixel values. The Number of Pixel Change Rate (NPCR) and the Unified Average Change in Intensity (UACI) are two parameters used to evaluate the sensitivity and randomness of an image. The algorithm is better, when the value of NPCR is close to 100 and UACI is just about 30. The NPCR and UACI are calculated by the equations (2) and (3), respectively.

\[
\text{NPCR} = \left( \frac{\sum_{i,j} |D(i,j) - 1|}{W \times H} \right) \times 100\% 
\]

\[
\text{UACI} = \left( \frac{1}{W \times H} \sum_{i,j} \left( \frac{|C1(i,j)-C2(i,j)|}{255} \right) \right) \times 100\% 
\]

Where,
- C1 and C2 are the cipher images whose plain images have a slight difference.
- C1(i, j) denotes the pixel position in C1, and C2(i, j) denotes the pixel position in C2.
- D(i, j) is determined by C1(i, j) and C2(i, j), namely if C1(i, j) = C2(i, j) then D(i, j)=1; otherwise, D(i, j)=0.
- W and H are width and height of the image.

The resulting NPCR and UACI values are tabulated in Table V to Table VIII for both existing and proposed methods.

<table>
<thead>
<tr>
<th>Table V NPCR VALUES OF PROPOSED METHOD (in %)</th>
</tr>
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<tbody>
<tr>
<td>Case</td>
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<tr>
<td>------</td>
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</table>

<table>
<thead>
<tr>
<th>Table VI NPCR VALUES OF EXISTING METHODS (in %)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Existing Methods</td>
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<tr>
<td>Ref. [12]</td>
</tr>
<tr>
<td>Ref. [23]</td>
</tr>
<tr>
<td>Ref. [25]</td>
</tr>
<tr>
<td>Ref. [26]</td>
</tr>
<tr>
<td>Ref. [27]</td>
</tr>
</tbody>
</table>

The experimental NPCR value of proposed method is similar to the existing methods suggested in [12], [26], [27], and approximately close to the methods proposed in [23], [24] and [25].

<table>
<thead>
<tr>
<th>Table VII UACI VALUES OF PROPOSED METHOD (in %)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Cipher Images</td>
</tr>
<tr>
<td>C1 Vs C2</td>
</tr>
<tr>
<td>C1 Vs C3</td>
</tr>
<tr>
<td>C1 Vs C4</td>
</tr>
<tr>
<td>C2 Vs C3</td>
</tr>
<tr>
<td>C2 Vs C4</td>
</tr>
<tr>
<td>C3 Vs C4</td>
</tr>
</tbody>
</table>

Where, C1, C2, C3 and C4 are encrypted Lena images using baboon, peppers, cameraman, and coin as key images, respectively.

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TABLE VIII

<table>
<thead>
<tr>
<th>Existing Methods</th>
<th>UACI Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Ref. [12]</td>
<td>28.6210</td>
</tr>
<tr>
<td>Ref. [23]</td>
<td>32.069</td>
</tr>
<tr>
<td>Ref. [24]</td>
<td>33.58</td>
</tr>
<tr>
<td>Ref. [25]</td>
<td>32.821</td>
</tr>
<tr>
<td>Ref. [26]</td>
<td>&gt; 24.94</td>
</tr>
<tr>
<td>Ref. [27]</td>
<td>&gt; 33.4</td>
</tr>
</tbody>
</table>

The measured UACI value of the proposed method is approximately close to the existing methods [25], [23], [24], and [27], and better than the methods suggested in [26], and [12].

VI. CONCLUSION

In this work, a novel encryption method is applied on digital images and found to give significantly improved performance. A comparable performance improvement is obtained in histogram, correlation, key space, and encryption speed. The histograms of the encrypted images are flat and so it is difficult to identify images by analyzing the histograms. The cross correlation value indicates that there is a weak relationship between the original and encrypted images. The correlation between adjacent pixels in the encrypted image is approximately close to zero. The cut test and dispersion test results show that the proposed method provides acceptable pixel permutation. The proposed method has approximately 10^{509} key space for an image of size 256 x 256 pixels to resist the brute force attack. The real time implementation choices confirm that the proposed method consumes less encryption time and could be robust for real time applications. The values of NPCR and UACI are greater than 99.4% and 30%, respectively, and thus the proposed method is resistant to differential attack.

REFERENCES


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