# Matching and Merging of Ontologies Using Conceptual Graphs

Gopinath Ganapathy, Ravi Lourdusamy

Abstract-Knowledge management applications need to determine whether two or more knowledge representations encode the same knowledge. Solving this matching problem is hard because representations may encode the same content but differ substantially in form. Previous approaches to this problem have used either syntactic measure or semantic knowledge to determine the distance between two representations. The aim of this article is to define matching of m conceptual graphs and present the mathematical aspects of matching and binding of m conceptual graphs with fitness. The algorithms developed for matching and merging of *m* conceptual graphs are illustrated.

#### Index Terms-Binding, Taxonomy, Transformation, Fitness.

#### I. INTRODUCTION

A requirement common to many knowledge applications is to determine whether two or many knowledge representations, encoded using the same ontology, capture the same knowledge. The task of determining whether two or more representations encode the same knowledge is treated as a graph matching problem. The knowledge representation is encoded using conceptual graph. The representations capture the same knowledge if their corresponding graphs match. The multiple encoding of the same knowledge rarely match exactly, so a matcher must be flexible to avoid a high rate of false-negatives. However, a matcher that is too flexible can suffer from a high rate of false-positives [6]. This problem has various causes, including

- (i) The ontology is expressible enough to allow the same information to be encoded in different ways
- (ii) The representations are built by different knowledge engineers (or computer programs), raising the likelihood they differ
- (iii) The representations are large, increasing the opportunity for differences.

Previous solutions to this problem have produced two types of matchers. Syntactic matchers use only the graphical form of the representations, judging their similarity by the amount of common structures shared [1], [2] or the number of edit operations required to transform one graph into the other [4], [7], [8], [10]. Approaches that focus on the amount of shared common structures do not handle mismatches. Approaches that use edit operations can handle mismatches but are sensitive to the cost assigned to the edit operations and tuning these parameters optimally is problematic.

In contrast, semantic matchers use knowledge, stored in an ontology, of the terms referenced in the representations. Semantic matchers use this knowledge to determine the match

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naries which will be useful for further discussion. Definition 2.1: [6] Transitive and part ascendant transformations conform to a more general notion called 'transfers through'. A relation r transfers through another relation r'if  $X \xrightarrow{r} V \xrightarrow{r'} Z \longrightarrow X \xrightarrow{r} Z$ 

Definition 2.2: [6] A triple is a 3-tuple of the form (head, relation, tail) where head and tail are concepts or instances (i.e., nodes in a conceptual graph) and relation is an edge in the graph. Every two nodes connected by an edge in a conceptual graph can be mechanically converted into a triple and hence a conceptual graph into a set of triples.

of two representations. [3], [5], [6], [12]. The knowledge

encoded can be equivalently encoded as conceptual graphs

[9]. The algorithm for matching of two graphs is discussed

in [11]. In this paper the generalized algorithm for matching

of m conceptual graphs is developed. The fitness of the

matched graph is further investigated. Further the algorithm

The article is organized as follows: Section 1 deals with

an introduction. The terminologies involved are outlined in

Section 2. In Section 3, the algorithm for matching mconceptual graphs is illustrated. The algorithm for merging

m- conceptual graphs is illustrated in Section 4 and in

**II. PRELIMINARIES** 

In this section, we present some definitions and prelimi-

is extended to merge m-conceptual graphs.

Section 5 conclusion is presented.

Definition 2.3: The *n* triples  $t_1 = (head_1, relation_1, relation_1)$  $tail_1$ ,  $t_2 = (head_2, relation_2, tail_2), \ldots, t_n = (head_n, tail_n), t_n$  $relation_n, tail_n$ ) of graph G align if  $head_1 \geq head_2 \geq$  $\dots \geq head_n, uneq(relation_1) \geq uneq(relation_2) \geq \dots \geq$  $uneg(relation_n)$  and  $tail_1 \geq tail_2 \geq ... \geq tail_n$ . The uneg(relation) unnegates relation if it is negated otherwise returns the relation.

Definition 2.4: For  $\ell = \{(t_{11}, t_{21}, \dots, t_{m1}), (t_{12}, t_{22}, \dots, t_{m2}), \dots, (t_{1n_1}, t_{2n_2}, \dots, t_{mn_m})\}$ , a list of aligned triples of m graphs, the bindings for  $\ell$  i.e.,

$$b(\ell) = \{(head_{11}/head_{21}/.../head_{m1}, tail_{11}/tail_{21}/.../tail_{m1}), \\ (head_{12}/head_{22}/.../head_{m2}, tail_{12}$$

$$\begin{aligned} tail_{22}/.../tail_{m2}), ..., (head_{1n_1}/head_{2n_2}/.../head_{mn_m}, \\ tail_{1n_1}/tail_{2n_2}/.../tail_{mn_m}) \end{aligned}\}. \end{aligned}$$

#### III. MATCHING OF m- CONCEPTUAL GRAPHS

## A. Binding of m- conceptual Graphs

Given m graphs  $G_1 = \{t_{11}, t_{12}, ... t_{1n_1}\}, G_2 =$  $\{t_{21}, t_{22}, ..., t_{2n_2}\}, ..., G_m = \{t_{m1}, t_{m2}, ..., t_{mn_m}\}$  where  $n_1, n_2, ..., n_m$  are the number of triples of  $G_1, G_2, ..., G_m$ respectively and a set of r transformations R where R = $\{R_1, R_2, ..., R_r\}$ . The aim of the algorithm is to find

(1)

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a common subgraph of  $G_1$ ,  $G_2$ ,... and  $G_m$  called SG. Construct a list M of all possible alignments between the triples of  $G_1$ ,  $G_2$ ,... and  $G_m$ . Each element of Mis of the form  $\ell = \{(t_{11}, t_{21}, ..., t_{m1}), (t_{12}, t_{22}, ..., t_{m2}), ...(t_{1n_1}, t_{2n_2}, ..., t_{mn_m})\}$ . The generalized algorithm for finding a match between m representations, is presented as Algorithm-1. The steps for finding a match between mrepresentations is illustrated with an example of three graphs,  $G_1$ ,  $G_2$  and  $G_3$  generated by organization structure of three hospitals shown in Figures – 1, 2 and 3. For reference, we label each triple in  $G_1$  with a unique number from 1 to 22, each triple in  $G_2$  with a unique lowercase letter from A to V and each triple in  $G_3$  with a unique lowercase letter from a to z and from aa to hh. We use subscripts to differentiate terms that appear multiple times (e.g., Hospital\_Model\_1).



Figure - 1 Graph  $G_1$ 

Algorithm - 1: Outline of the generalized matching algorithm

1) M = NIL and  $\ell = NIL$  FOR each triple  $t_{1i_1}$  in  $G_1$ FOR each triple  $t_{2i_2}$  in G2

FOR each triple 
$$t_{mi_m}$$
 in  $G_m$   
If  $t_{1i_1}, t_{2i_2}, ..., t_{mi_m}$  are aligned  
THEN add  $(t_{1i_1}, t_{2i_2}, ..., t_{mi_m})$  to  $\ell$ .  
ADD  $\ell$  to  $M$  and reset  $\ell$  to NIL.

- 2) Use *M* to construct a common subgraph of  $G_1, G_2, ..., G_m$  called *SG*.  $SG = \{(t_1, \dots, t_{n-1}), (t_{n-1}, \dots, t_{n-1})\}$ 
  - $SG = \{(t_{11}, ..., t_{m1}), (t_{12}, ..., t_{m2}), ..., (t_{1n_1}, ..., t_{mn_m})\}$ where  $(t_{1i}, ..., t_{mi})$  are the aligned triples of  $G_1, G_2, ..., G_m$ respectively.
- 3) IF SG is inconsistent THEN stop and return NIL

4) FOR each rule R<sub>i</sub> in R, FOR each j = 1, 2, ..., m FOR each k = 1, 2, ..., m IF R<sub>i</sub> is applicable to G<sub>j</sub> with respect to G<sub>k</sub>

- THEN apply  $(R_i, G_j, G_k)$ 5) FOR each j = 1, 2, ..., m
- For each unaligned triple  $t_{jij}$  in  $G_j$   $IF t_{1i_{1j}}, t_{2i_{2j}}, ..., t_{mim_j}$  are aligned and  $b(\{(t_{1i_{1j}}, t_{2i_{2j}}, ..., t_{mim_j})\})$  is consistent with b(SG),  $THEN \text{ add}(t_{1i_{1j}}, t_{2i_{2j}}, ..., t_{mim_j})$  to SG and break. UNTIL SG reaches quiescence go to step 4.
- 6) RETURN SG



In step 1 the generalized algorithm compares each triple in  $G_1$  with each triple in  $G_2$ ,  $G_3$ ,..., $G_m$  to find all possible alignments. In our example  $t_{11}$  aligns with  $t_{21}$  and  $t_{31}$ . Triple  $t_{11}$ , however, does not align with triple  $t_{22}$  and its combinations because the relations differ. The initial match is denoted by M.  $M = \{\ell_1, \ell_2, \dots, \ell_p\}$  where p is the total number of possible alignments. Each element of Mis a list called  $\ell_i$ . For example  $\{(t_{11}, t_{21}, t_{31})\}$  is called  $\ell_1$ ,  $\{(t_{12}, t_{22}, t_{32})\}$  is called  $\ell_2$ , etc. In step 2 the generalized algorithm uses M to construct common subgraph of  $G_1$ ,  $G_2,...,G_m$  called SG. The generalized algorithm begins by selecting a member,  $\ell_i$ , of M to serve as the seed of the construction process (recall that  $M = \{\ell_1, \ell_2, ..., \ell_s\}$ ) and  $\ell_i = \{(t_{1i_{11}}, t_{2i_{21}}, t_{mi_{m1}}), (t_{1i_{12}}, t_{2i_{22}}, t_{mi_{m2}}), ..., \}$  $(t_{1i_{1k_i}}, t_{2i_{2k_i}}, t_{mi_{mk_i}})$ }, i = 1, 2, ..., s. This seed is selected based on a heuristic scoring function

$$h(\ell_i) = \frac{1}{k_i} \left\{ \sum_{j=1}^{k_i} n(head_{1i_{1j}}/head_{2i_{2j}}/.../head_{mi_{mj}}) + n(tail_{1i_{1j}}/tail_{2i_{2j}}/..../tail_{mi_{mj}}) \right\}$$
(2)

where  $head_{1i_{1j}}/head_{2i_{2j}}/, ..., /head_{mi_{mj}}$  and  $tail_{1i_{1j}}/tail_{2i_{2j}}/, ..., /tail_{mi_{mj}}$  are the bindings of  $t_{1i_{1j}}, t_{2i_{2j}}, ..., t_{mi_{mj}}$  and n(b) is the number of times the bindings b occurs in binding(M). This function h is a heuristic that favors those  $\ell_i$  in M with interconnectivity. Bindings that occur frequently indicate

high interconnectivity. We want to select these  $\ell_i$  as the seeds because they have more potential for allowing larger common subgraphs to be constructed. Therefore, the algorithm selects the  $\ell_i$  in M with the highest score, as determined by the function h. SG is extended with those pairs of aligned triples in M whose bindings intersect the binding of the pairs in SG. Pairs in M that extend SG are removed from M along with  $\ell_j$  they belong to. This process is repeated until SG can no longer be extended.

In steps 3-5, the generalized algorithm checks if SGis consistent. SG is inconsistent if it contains an aligned *m*-tuples of triples  $(t_{1i}, t_{2i}, ..., t_{mi})$  where the relation of at least one  $t_{ji}$  is negated and the relation of at least one  $t_{ki}$  is not negated. If SG contains such a m-tuples, then generalized algorithm stops and returns NIL. Otherwise, the generalized algorithm applies transformations to improve the match(i.e., steps 4 and 5). Steps 4 and 5 are repeated until SG reaches quiescence. In step 4, the generalized algorithm applies transformations to resolve mismatches among  $G_1, G_2, ..., G_m$ . In step 5, the generalized algorithm will try to align additional triples among  $G_1, G_2, ..., G_m$ . Step 5 is like step 1 except that it focuses on the unaligned triples. We have identified a set of transformations for the health care domain. These transformations are used to improve the matching of *m*-conceptual graphs in the domain. Returning to our example, the triples of  $G_1$ ,  $G_2$ and  $G_3$  are as follows:

 $G_1 = \{(Hospital_Model_1, hasAct, Allot_Bed_1), (Hospital_Model_1, Model_1), (Hospital_Model_1), (Hospit$ hasAct, New\_Renew\_Register\_patient\_1), (Hospital\_Model\_1, hasAct, Admit\_patient\_1), (Hospital\_Model\_1, hasParticipation, Registration\_1), (Hospital\_Model\_1, hasAct, Diagnose\_patient\_1), (Hospital\_Model\_1, hasRole, Patients attendee 1), (Admit patient 1, for, Registration 1), (Registration\_1, has, Admit\_patient\_1), (Registration\_1, has, Allot\_Bed\_1), (Allot\_Bed\_1, for, Registration\_1) (New\_Renew\_Register\_patient\_1, for, Registration\_1), (Registration\_1, has, New\_Renew\_Register\_patient\_1), (Registration\_1, has, Diagnose\_patient\_1), (Diagnose\_patient\_1, for, Registration\_1), (Patients\_attendee\_1, hasAsParticipant, Registration\_1), (Registration\_1, participatesIn, Patients\_attendee\_1), (Hospital\_Model\_1, hasRole, Nurse\_1), (Hospital\_Model\_1, hasEntity, Stethoscope\_1), isPlayedBy, Physician\_1), (Physician\_1, plays, (Stethoscope\_1, (Nurse\_1, hasAsParticipant, Registration\_1), Stethoscope 1), (Registration\_1, participatesIn, Nurse\_1)}

 $G_2 = \{ (Hospital_Model_2, hasAct, Allot_Bed_2), (Hospital_Bed_2), (Hospital_Be$ hasAct, New\_Renew\_ Registration\_ patient\_2), (Hospital\_Model\_2, hasAct, Admit\_patient\_2), (Hospital\_Model\_2, hasParticipation, Registration\_2), (Hospital\_Model\_2, hasAct, Diagnose\_patient\_2), (Hospital\_Model\_2, hasRole, Patients\_attendee\_2), (Admit\_patient\_2, for, Registration\_2), (Registration\_2, has, Admit\_patient\_2), (Registration\_2, has, Allot\_Bed\_2), (Allot\_Bed\_2, for, Registration\_2), (New\_Renew\_Register\_patient\_2, for, Registration\_2), (Registration\_2, has, New\_Renew\_Register\_patient\_2), (Registration\_2, has, Diagnose\_patient\_2), (Diagnose\_patient\_2, for. Registration\_2), (Patients\_attendee\_2, hasAsParticipant, Registration\_2), (Registration\_2, participatesIn, Patients\_attendee\_2), (Hospital\_Model\_2, hasRole, Registration\_staff\_2), (Hospital\_Model\_2, Check\_weight\_and\_BP\_of\_patient\_2), has Act. (Check\_weight\_and\_BP\_of\_patient\_2, for, Registration\_2), (Registration\_2, has, Check\_weight\_and\_BP\_of\_patient\_2), (Registration\_Staff\_2, hasAsParticipant, Registration\_2), (Registration\_2, participatesIn, Registration\_Staff\_2)

 $G_3 = \{(Hospital_Model_3, hasAct, Allot_Bed_3), (Hospital_Model_3, Model_3, Model_$ 

hasAct. New\_Renew\_Register\_patient\_3), Hospital Model 3. hasAct. Admit\_patient\_3, (Hospital\_Model\_3, hasParticipation, Registration\_3), (Hospital\_Model\_3, hasAct. Diagnose patient 3). (Hospital\_Model\_3, hasRole, Patients\_attendee\_3), (Admit\_patient\_3, Registration\_3), (Registration\_3, Admit\_patient\_3), for. has, (Registration\_3, has, Allot\_Bed\_3), (Allot\_Bed\_3, for, Registration\_3), (New\_Renew\_Register\_ patient\_3, for, Registration\_3),(Registration\_3, New\_Renew\_Register\_patient\_3), (Registration\_3, has. has, Diagnose\_patient\_3), (Diagnose\_patient\_3, for, Registration\_3), (Patients\_attendee\_3, hasAsParticipant, Registration\_3 ), (Registration\_3, ParticipatesIn, Patients\_attendee\_3), (Hospital \_Model\_3, hasRole, Billing\_Assistant\_3), (Hospital\_Model\_3, hasEntity, Stethoscope\_3), (Hospital\_Model\_3, hasEntity, BP\_Apparatus\_3), (Hospital\_Model\_3, hasRole, Nurse\_3), (Hospital\_Model\_3, hasAct, Check\_weight\_and\_BP\_ of\_patient\_3), (Hospital\_Model\_3, hasRole, Registration staff 3). (Nurse\_3, hasAsParticipant, Registration\_3), (Registration\_3, ParticipatesIn, Nurse\_3), (Registration\_3, has, Check\_weight\_ and\_BP\_of\_patient\_3), (Check\_weight\_and\_BP\_of\_patient\_3, for, Registration\_3), (Physician\_3, plays, Stethoscope\_3, (Stethoscope\_3, isPlayedBy, Physician\_3), (Collect\_money\_for \_consolidated\_bill\_3, for, Billing\_3), (Billing\_3, Collect\_money\_for\_consolidated\_bill\_3), (Billing\_Assistant\_3, has, has. Billing\_3), (Billing\_3, ParticipatesIn, Billing\_assistant\_3), (Prepare\_consolidated\_bill\_3, for, Billing\_3), (Billing\_3, has, Prepare consolidated bill 3)

In step-1, the algorithm compares each triple in  $G_1$  with each triple in  $G_2$  and  $G_3$  to find all possible alignments. Triple 1 of  $G_1$  aligns with A of  $G_2$  and a of  $G_3$ , Triple 2 of  $G_1$  aligns with B of  $G_2$  and b of  $G_3$  and so on. Hence the matched triples of  $G_1, G_2$  and  $G_3$  are as follows: Binding(M) = {(1, A, a), (2, B, b), (3, C, c), (4, D, d), (5, E, e), (6, F, f), (7, G, g), (8, H, h), (9, I, i), (10, J, j), (11, K, k), (12, L, l), (13, M, m), (14, N, n), (15, O, o), (16, P, p)} = { $\ell_1, \ell_2, \ell_3, \ell_4, \ell_5, \ell_6, \ell_7, \ell_8, \ell_9, \ell_{10}, \ell_{11}, \ell_{12}, \ell_{13}, \ell_{14}, \ell_{15}, \ell_{16}$ }.

In step-2, hscore is calculated as follows:  $hscore(\ell_1) = hscore(\ell_2) = hscore(\ell_3) = 10, hscore(\ell_4) = 18, hscore(\ell_5) = hscore(\ell_6) = 10, hscore(\ell_j) = 14$ , where j = 7, 8, ..., 16. Select  $\ell_4$  and remove it from M and hence the subgraph  $SG = \{(4, D, d)\}.$ 

Binding(M)= { $\ell_1$ ,  $\ell_2$ ,  $\ell_3$ ,  $\ell_5$ ,  $\ell_6$ ,  $\ell_7$ ,  $\ell_8$ ,  $\ell_9$ ,  $\ell_{10}$ ,  $\ell_{11}$ ,  $\ell_{12}$ ,  $\ell_{13}$ ,  $\ell_{14}$ ,  $\ell_{15}$ ,  $\ell_{16}$ }. The head of  $\ell_4$  is (Hospital\_Model\_1/ Hospital\_Model\_2/Hospital\_Model\_3) is intersecting with  $\ell_1$ ,  $\ell_2$ ,  $\ell_3$ ,  $\ell_5$ ,  $\ell_6$ . Remove all from M and hence the subgraph of SG {(4, D, d), (1, A, a), (2, B, b), (3, C, c), (5, E, e), (6, F, f)}. Binding(M) = { $\ell_7$ ,  $\ell_8$ ,  $\ell_9$ ,  $\ell_{10}$ ,  $\ell_{11}$ ,  $\ell_{12}$ ,  $\ell_{13}$ ,  $\ell_{14}$ ,  $\ell_{15}$ ,  $\ell_{16}$ }. The tail of  $\ell_4$ is (Registration\_1/Registration\_2/Registration\_3) is intersecting with  $\ell_7$ ,  $\ell_8$ ,  $\ell_9$ ,  $\ell_{10}$ ,  $\ell_{11}$ ,  $\ell_{12}$ ,  $\ell_{13}$ ,  $\ell_{14}$ ,  $\ell_{15}$ ,  $\ell_{16}$ . Hence the subgraph

$$\begin{split} SG =& \{(4, D, d), (1, A, a), (2, B, b), (3, C, c), \\ & (5, E, e), (6, F, f), (7, G, g), (8, H, h), \\ & (9, I, i), (10, J, j), (11, K, k), (12, L, l), \\ & (13, M, m), (14, N, n), (15, O, o), (16, P, p)\} \end{split}$$

The result is the maximal subgraph of  $G_1$ ,  $G_2$  and  $G_3$  which is shown in Figure 4.



Fig 4 Matched Graph





Figure - 5 Transformation

SG is consistent, hence we apply transformations to improve the match. For example transformation  $R_1$  is applied to graphs  $G_1, G_2$  and  $G_3$  and hence a triple is added to graphs  $G_1, G_2$  and  $G_3$ . Because of the added triple another match is identified and that is added in SG. The added triple is shown as an edge of dotted line in Figure 4.

#### B. Fitness of the binding subgraph

In this section we present the fitness of the matched Graph. SG is returned along with a numeric score reflecting the fitness of the match between  $G_1, G_2, ..., G_m$ . This score is used in situations where one graph is matched with a set of graphs to select the best match. To calculate the fitness score, we apply equation (3) to SG where

$$SG = \{(t_{1j_1}, t_{2j_1}, \dots, t_{mj_1}), \\ (t_{1j_2}, t_{2j_2}, \dots, t_{mj_2}), \dots, \\ (t_{1j_k}, t_{2j_k}, \dots, t_{mj_k})\}$$

and the *m*-tuples  $(t_{1j_i}, t_{2j_i}, ..., t_{mj_i})$ , for j = 1, 2, ..., k are aligned triples of  $G_1, G_2, ..., G_m$  respectively.

$$Fitness(SG) = \frac{\sum_{i=1}^{k} score(t_{1j_i}, t_{2j_i}, ..., t_{mj_i})}{min\{(|G_1| + t_{r1}), (|G_2| + t_{r2}), ..., (|G_m| + t_{rm})\}}$$
(3)

where  $|G_s|$  is the number of triples in  $G_s$  and  $t_{rs}$  is the number of additional triples added to  $G_s$ , for  $s = 1 \cdots m$ , by applying transformations. The score for each *m*-tuple of triples  $(t_{1j_i}, t_{2j_i}, ..., t_{mj_i})$  is computed using a function score as given by equation (4).

$$score(t_{1j_i}, t_{2j_i}, ..., t_{mj_i}) = \frac{\frac{1}{x+1} + \frac{1}{y+1} + \frac{1}{z+1}}{3}$$
 (4)

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where,  $x = taxdist(head_{1j_i}, head_{2j_i}, ..., head_{mj_i})$ ,  $y = taxdist(rela_{1j_i}, rela_{2j_i}, ..., rela_{mj_i})$ ,  $z = taxdist(tail_{1j_i}, tail_{2j_i}, ..., tail_{mj_i})$  and

$$taxdist(c_1, c_2, ..., c_m) = \sum_{i=1}^{m-1} taxdist(c_i, c_{i+1})$$

the function  $taxdist(c_i, c_{i+1})$  is the taxonomic (or semantic) distance [9] between two concepts  $c_i$  and  $c_{i+1}$ . We calculate the taxonomic distance between two concepts as the minimum number of taxonomic edges that needs to be traversed to reach  $c_i$  from  $c_{i+1}$ . After transformations have been applied, SG is returned along with a numeric score reflecting the fitness of the match between  $G_1, G_2,..., G_m$ . This score is computed based on the number of matched triples over the size of the graphs being matched [11]. The fitness score is also computed using a simple formula as given by equation (5)

$$Fitness(SG) = \frac{No. of Triples of SG}{Maximum of n_1, n_2, ..., n_m}$$
(5)

In our example, the fitness of SG is approximately 0.5.

# IV. MERGING OF m-conceptual graphs

In this section we present an algorithm to merge mconceptual graphs after determining the matched graph. Let  $p_1, p_2, ..., p_m$  be the number of unmatched triples of the graphs  $G_1, G_2, ..., G_m$ . Assume that the matched graph SG as G and consider each unmatched triple  $t_{jij}$  of a graph  $G_i$ . Search G to determine the node containing the head of  $t_{jij}$ . After determining such node,  $t_{ji}$  is added as a subnode of the searched graphs node. After adding all the unmatched triples in G, the resultant graph is the merged graph of m Graphs  $G_1, G_2, \ldots, G_m$ . Returning to our example the merged graph of  $G_1, G_2$ , and  $G_3$  is shown in Fig 6.



Legend: H\_M: Hospital\_Model, Reg: Registration, p: patient, N\_R: New\_Renew

Algorithm - 2: Outline of the generalized merging algorithm

(1) Set SG as G

(2) FOR each unmatched triple  $t_{1i_1}$  in  $G_1$ FOR each unmatched triple  $t_{2i_2}$  in  $G_2$ 

FOR each unmatched triple  $t_{mi_m}$  in  $G_m$  search G for the head of  $t_{ji_j}$ , in  $n_i$ 

Determine the node where there is a match add  $t_{ji_j}$  as a subnode for that searched node and return the resultant graph as G.

(3) RETURN G

### V. CONCLUSION

Algorithms for matching and merging of m-conceptual graphs are developed and the theoretical aspects are discussed in detail. Further these algorithms can be implemented using any ontology management software to support its application to the health care domain.

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