The PDBO Algorithm for Discrete Time, Cost and Quality Trade-off in Software Projects with Task Quality Estimated by COQUAMO

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Abstract—Software project manager's goal is getting an optimal allocation of time, cost, and quality of each task/activity in the software project such that total time and cost of the project is minimized while project quality is maximized. Accordingly, mathematical and meta-heuristic algorithms have been developed in order to solve such riddle. This research paper introduces a new proposed meta-heuristic, problem data based optimization (PDBO), algorithm to solve the discrete time-cost-quality trade-off problem (DTCQTP) in software projects. PDBO decides the preferred modes of performing tasks where task quality in each mode is estimated using Constructive quality model (COQUAMO) based on task cost in that mode. An example is given at the end to show the trade-off analysis between project's cost, time and quality.

Index Terms—Meta-heuristic algorithms, Time cost quality trade-off, Quality estimation, Software projects

INTRODUCTION

One main task for project managers is to administrate projects under concern and achieve the required goals within the plan. Improving the resources allocation to guarantee minimum cost, time and high quality is an obligatory task for such administration [8, 14]. Accordingly, many researchers have devoted much effort to solve such riddle, on one hand, some of these researches considered continues mode for the time, cost and quality [13]. On the other hand, multiple modes for each activity depending on discrete models have been considered [10]; in such situations there are alternative approaches for completing each task/activity, each having its own time, cost, and quality considerations. Differences in quality can arise due to bids offered by competing contractors to complete specific tasks. Even different bids by the same subcontractor could imply different quality levels. For example, subcontractors might have some flexibility with time and cost that would result in different quality levels for the same task. This can also be true for alternative work plans offered in house. For example, in analysis task there are choices (techniques) such as interview, questionnaire, observation, documentation review, etc related to gather facts about the software been developed.

Each of the possible alternatives will achieve different levels of time, cost, and quality associated with this task. Accordingly, mathematical and meta-heuristic techniques are taken into account to solve such problems [1, 2, 3, 5, 8, 15]. In the two cases, continues mode and discrete mode, task quality is measured based on managers' judgment which is expressed by values such as 90%, 80%, etc or by quality indicators which don't reflect exactly number of defects in a task. Accordingly, in this paper, the task in specific mode, its quality is estimated by COQUAMO based on its cost in that mode and thus a total number of residual defects for all tasks in the chosen modes reflect the quality of the software project quantitatively. To solve the problem in this discrete case, discrete time-cost-quality trade-off problem (DTCQTP), PDBO algorithm is used. The paper also introduces an example that shows the trade-off analysis between project's time, cost and quality. The rest of this paper is organized as follows: in Section II related work, section III problem definition, section IV COQUAMO, section V software project time, cost and quality, section VI methodology, section VII example and section VIII concludes and discusses the paper.

I. RELATED WORK

Discrete time-cost-quality trade-off problem DTCQTP is the problem of optimizing time, cost and quality based on discrete mathematical models; it is an extension of discrete time-cost tradeoff problem (DTCTP) by taking quality into consideration [1]. The first work initiated by Babu and Suresh in 1996 who claimed that the cost and quality of a completed project may be affected by activity crashing. Thus they developed a solution procedure based on linear models which optimizes time, cost and quality in continuous mode [15]. Later on in 1999 this procedure has been applied on real cement factory construction project in Thailand for evaluation. The practicality and feasibility of the time and cost optimization models were demonstrated. But the methodology of the quality measurement were said to be over subjective and inaccurate [13]. Although this procedure can assist managers in making tradeoff decisions by providing valuable information, it disregards the multiple modes for different activities [15]. In fact, the introduction of quality is of great importance to the research of the time-cost-quality trade-off problem, however it has to be pointed out that there relatively little accurate methods to quantify quality. Moreover, despite that the linear relations between time and cost, time and quality are accepted by some scholars in their research, these assumptions are not practical in real life projects, especially when the existence
of budget threshold in the time-cost curve is taken into consideration.

After that, several works considered discrete models by many authors and many algorithms have been developed to solve them. El-Rayes et al in 2005 initially formulated the first discrete time-cost-quality trade-off model. Meanwhile Genetic Algorithm (GA) was employed to search for the Pareto optimal solutions, which provided new visions to solve the models [21]. Tareghian et al in 2006 developed three inter related binary integer programming models for DTCQTP such that each optimizes one of the given three entities by assigning desired bounds on the other two and used lingo software for optimization. However, they proposed to use hybrid meta-heuristic combining scatter search and electromagnetism ideas to solve large scale DTCQTP in future work due to increasing number of iterations of the algorithm that LINGO optimizer uses [4]. Afshar et al in 2007 developed a new meta heuristic, multi-colony ant algorithm, for optimizing time-cost-quality tradeoff to generate optimal/near optimal solutions [8]. Because DTCQTP is NP-Hard, Iranmanesh et al in 2008 proposed a meta-heuristic based on GA to solve such problem [7]. Refaat et al in 2010 developed a practical software system using a multi-objective genetic algorithm (GA) for optimizing time-cost-quality tradeoff simultaneously to help planers in decision making [5] and Shankar et al in 2011 analyzed project scheduling problem in terms of time, cost and quality [6]. All of these mentioned works, task quality is measured based on manager’s judgment or using quality indicators. Shahsavari et al in 2010 developed a mathematical model for discrete time, cost and quality tradeoff problem using a novel hybrid genetic algorithm (NHGA) [2]. Moreover, in 2012 in order to handle project quality uncertainty, NHGA has been applied associated with fuzzy logic by assuming time and cost as crisp variables, while quality as linguistic variable [3]. Roya et al 2013 estimated task quality based on its time and cost using fuzzy logic; however project, personnel, platform and product attributes are not taken into account [14]. The existing literature provides a broad vision for research of the time-cost-quality trade-off problem; where non-trivial defects are reduced by using more resources and overcomes the deadline problem. However, more resources lead to cost increasing. Recently, project managers’ main consideration is to improve the project quality while reducing both the time needed and the cost leading to discrete time, cost, quality trade-off problem (DTCQTP) [1-8]. Accordingly, many meta heuristic algorithms have been devoted to solve such problem such as genetic [2, 3], practical swarm optimization (PSO) [1] and Multi-colony ant optimization [8] algorithms. DTCQTP has multiple efficient solutions, but in this work, a single solution is obtained in terms of minimum cost and time with maximum quality.

III. CONSTRUCTIVE QUALITY MODEL (COQUAMO)

Constructive quality model (COQUAMO) is an extension of the existing constructive cost model (COCOMO II) and consists of two sub-models, defects introduction sub-model DI and defects removal sub-model DR, as in fig 1 and fig 2 respectively.

The DI sub-model’s inputs include source lines of code and/or function points as the sizing parameter, adjusted for both reuse and breakage, and a set of 21 multiplicative DI-drivers divided into four categories, platform, product, personnel and project. These 21 DI-drivers are a subset of the 22 cost parameters required as input for COCOMO II. Development flexibility FLEX driver has no effect on defect introduction and thus here its values for rating are set to 1.

The decision to use these drivers was taken after the author did an extensive literature search and did some behavioral analyses on factors affecting defect introduction. The outputs of DI sub-model are predicted number of non-trivial defects of requirements, design and code introduced during development life cycle; where non-trivial defects include:

- Critical (causes a system crash or unrecoverable data loss or jeopardizes personnel)
- High (causes impairment of critical system functions and no workaround solution exists)
- Medium (causes impairment of critical system function, though a workaround solution does exist).

Based on expert-judgment, an initial set of values to each of ratings of the DI drivers that have an effect on the number of defects introduced and overall software quality were proposed and we are used them in our implementation.

The aim of the defect removal (DR) model is to estimate the number of defects removed by several defect removal activities namely automated analysis AUTA, people reviews PEER and execution testing and tools EXTT. The DR model is a post-processor to the DI model. Each of these three defect removal profiles removes a fraction of the requirements, design and coding defects introduced from DI model. Each profile has 6 levels of increasing defect removal capability, namely ‘Very Low’, ‘Low’, ‘Nominal’, ‘High’, ’Very High’ and ‘Extra High’ with ‘Very Low’ being the least effective and ‘Extra High’ being the most effective in defect removal.

To determine the defect removal fractions (DRF) associated with each of the six levels (i.e. very low, low, nominal, high, very high, extra high) of the three profiles (i.e. automated analysis, people reviews, execution testing and tools) for each of the three types of defect artifacts (i.e. requirements defects, design defects and code defects), the author conducted a 2-round Delphi and we used the values of DRF resulted from 2-round Delphi in our implementation.
The inputs of DR sub-model include software size in thousand source lines of code KSLOC and/or function points, defect removal profiles levels and number of non-trivial defects of requirements, design and code from DI model. For more details about COQUAMO see [17] [19].

IV. SOFTWARE PROJECT TIME, COST and QUALITY

Time: is time required to develop software. Cost includes hardware and software costs, travel and training costs, effort costs (the most dominant factor in most projects) and effort costs overheads; costs of building, heating, lighting, costs of networking and communications and costs of shared facilities (e.g. library, staff restaurant, etc.) [16]. These costs are classified as direct cost which vary during project development such as travel costs and indirect cost which remain constant during time unit such as lighting costs [2]. Quality has been used in different contexts and has different definitions [17] which mean different things to different people [18], but in this research, quality is defined as number of residual defects in the activity. The defect is defined as a divergent of actual results from desired results [17].

V. METHODOLOGY

The following subsections explain task quality estimation, DTCQTP representation and assumptions and mathematical modeling.

A. Task Quality Estimation

After estimating software size and cost and time for every task/activity in every mode by any estimation method for software which will be developed in house or if bids offered from subcontractors to perform specific tasks such that each bid represent one mode, task quality in every mode is estimated by COQUAMO based on its cost in that mode.

To estimate task quality in every mode by COQUAMO based on its cost in that mode, we used the rules shown in table I and assume project manager knows the range of cost levels such as task cost is very low VL when it is between 0 and 20$ and it is low L when it is between 20 and 205$ and it is medium M when it is between 205 and 400$ and it is high H when it is between 400 and 600$ and it is very high VH when it is between 600 and 1000$ and it is extraordinary high EH when it is between 1000 and 2000$.

Table I is built based on COCOMO II and COQUAMO models which show the influence of cost drivers (DI-drivers) levels and defects removal profiles levels on effort and hence on cost of the tasks and also on defects introduced to and removed from the tasks.

Row labeled R with yellow color in table I represent “if part” while all other rows numbered from 1 through 25 with white color represent ”than part” e.g. if task cost level in specific mode is very low VL then this means that column correspond to cost which is labeled VL applies for all drivers and removal profiles i.e. PREC=EX, RESL=EX and so on.

According to task cost level, drivers and removal profiles levels will be chosen and entered into COQUALMO along with software size to estimate task quality.

Requirements include feasibility study and analysis tasks, so we assume that feasibility study represent 1/3 of Requirements and thus its defects =1/3 of requirements defects and analysis represent 2/3 of Requirements and thus its defects =2/3 of requirements defects. According to Jones report [20], documentation defects =0.60 per function point fp, requirements defects=1 per fp, design defects=1.25 per fp and code defects=1.75 per fp. So documentation defects =0.60/1=60% of requirements defects or =0.60/1.25=48% of design defects or =0.60/1.75=34% of code defects.
Determine ID driver

Estimated introduced defects in design:

$$D_{req}^{req} \cdot A1. (\text{Size})^{\beta_1} \cdot \prod_{j=1}^{21} (DI - Driver) j_{\text{Req}}$$

Estimated removed defects in design:

$$D_{req}^{req} = C1. D_{req}^{req} \cdot \prod_{j=1}^{3} (1 - DRF) r_{\text{Req}}$$

Estimated residual defects in analysis:

$$\text{Resid}^{\text{ANA}} = (D_{req}^{req} - D_{req}^{req})^{2/3}$$

Estimated residual defects in requirements:

$$\text{Resid}^{\text{DRF}} = (D_{req}^{req} - D_{req}^{req})^{\alpha_1} \cdot \prod_{j=1}^{3} (1 - DRF) r_{\text{Req}}$$

Estimated residual defects in design:

$$D_{req}^{req} = (D_{req}^{req} - D_{req}^{req})^{\alpha_1} \cdot \prod_{j=1}^{3} (1 - DRF) r_{\text{Req}}$$

Estimated residual defects in Feasibility study:

$$\text{Resid}^{\text{FS}} = (D_{req}^{req} - D_{req}^{req})^{\alpha_1} \cdot \prod_{j=1}^{3} (1 - DRF) r_{\text{Req}}$$

Estimated residual defects in documentation:

$$\text{Resid}^{\text{DOC}} = (D_{req}^{req} - D_{req}^{req})^{\alpha_1} \cdot \prod_{j=1}^{3} (1 - DRF) r_{\text{Req}}$$

Estimated residual defects in code:

$$D_{req}^{req} = C3. D_{req}^{req} \cdot \prod_{j=1}^{3} (1 - DRF) r_{\text{Req}}$$

Estimated residual defects in documentation:

$$\text{Resid}^{\text{Code}} = (D_{req}^{req} - D_{req}^{req})^{\alpha_1} \cdot \prod_{j=1}^{3} (1 - DRF) r_{\text{Req}}$$

Estimated residual defects in code:

$$\text{Resid}^{\text{Code}} = (D_{req}^{req} - D_{req}^{req})^{\alpha_1} \cdot \prod_{j=1}^{3} (1 - DRF) r_{\text{Req}}$$

Size is the size of the software project measured in terms of KSLOC (thousands of source lines of code, function points or any other unit of size), here KSLOC is used as software size measure.

$$B1, B2$$ and $$B3$$ accounts for economies / diseconomies of scale and are initially set to 1 as in [17, 19].

$$r = 1$$ to 3 for each DR profile, namely automated analysis, people reviews, execution testing and tools.

The flowchart, fig 4, shows how task quality in each mode is estimated.

### Table I: If-then rules

<table>
<thead>
<tr>
<th>R</th>
<th>Task Cost</th>
<th>VL</th>
<th>L</th>
<th>N</th>
<th>H</th>
<th>VH</th>
<th>XH</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>then PREC</td>
<td>EX</td>
<td>VH</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>2</td>
<td>then RESL</td>
<td>EX</td>
<td>VH</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>3</td>
<td>then TEAM</td>
<td>EX</td>
<td>VH</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>4</td>
<td>then PMAT</td>
<td>EX</td>
<td>VH</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>5</td>
<td>then FLEX</td>
<td>EX</td>
<td>VH</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>6</td>
<td>then RELY</td>
<td>VL</td>
<td>L</td>
<td>N</td>
<td>H</td>
<td>VH</td>
<td>VH</td>
</tr>
<tr>
<td>7</td>
<td>then DATA</td>
<td>L</td>
<td>L</td>
<td>N</td>
<td>H</td>
<td>VH</td>
<td>VH</td>
</tr>
<tr>
<td>8</td>
<td>then DOCU</td>
<td>VL</td>
<td>L</td>
<td>N</td>
<td>H</td>
<td>VH</td>
<td>VH</td>
</tr>
<tr>
<td>9</td>
<td>then CPLX</td>
<td>VL</td>
<td>L</td>
<td>N</td>
<td>H</td>
<td>VH</td>
<td>XH</td>
</tr>
<tr>
<td>10</td>
<td>then RUSE</td>
<td>L</td>
<td>L</td>
<td>N</td>
<td>H</td>
<td>VH</td>
<td>XH</td>
</tr>
<tr>
<td>11</td>
<td>then TIME</td>
<td>N</td>
<td>N</td>
<td>N</td>
<td>H</td>
<td>VH</td>
<td>VH</td>
</tr>
<tr>
<td>12</td>
<td>then STOR</td>
<td>N</td>
<td>N</td>
<td>N</td>
<td>H</td>
<td>VH</td>
<td>XH</td>
</tr>
<tr>
<td>13</td>
<td>then PVOL</td>
<td>L</td>
<td>L</td>
<td>N</td>
<td>H</td>
<td>VH</td>
<td>VH</td>
</tr>
<tr>
<td>14</td>
<td>then ACAP</td>
<td>VH</td>
<td>H</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>15</td>
<td>then AEXP</td>
<td>VH</td>
<td>H</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>16</td>
<td>then PCAP</td>
<td>VH</td>
<td>H</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>17</td>
<td>then PEXP</td>
<td>VH</td>
<td>H</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>18</td>
<td>then LTEX</td>
<td>VH</td>
<td>H</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>19</td>
<td>then PCON</td>
<td>VH</td>
<td>H</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>20</td>
<td>then TOOL</td>
<td>VH</td>
<td>H</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>21</td>
<td>then SCED</td>
<td>VH</td>
<td>H</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>22</td>
<td>then SITE</td>
<td>EX</td>
<td>VH</td>
<td>N</td>
<td>L</td>
<td>VL</td>
<td>VL</td>
</tr>
<tr>
<td>23</td>
<td>then AUTA</td>
<td>VL</td>
<td>L</td>
<td>N</td>
<td>H</td>
<td>VH</td>
<td>XH</td>
</tr>
<tr>
<td>24</td>
<td>then PEER</td>
<td>VL</td>
<td>L</td>
<td>N</td>
<td>H</td>
<td>VH</td>
<td>XH</td>
</tr>
<tr>
<td>25</td>
<td>then EXTT</td>
<td>VL</td>
<td>L</td>
<td>N</td>
<td>H</td>
<td>VH</td>
<td>XH</td>
</tr>
</tbody>
</table>

B. DTCQTP Representation and Assumptions

PDBO algorithm assumes that the DTCQTP has the representation of activity-on-node network and also assumes that every node (activity) in a network has virtual edges to all its modes as in fig 5. Activities 1 and 5 have virtual edges to all their modes where M1 is the first mode and Mn is last mode of an activity and the other activities have virtual edges to their modes similar to activities 1 and 5. Activities S and F don not have modes and therefore do not have virtual edges to modes because these are dummy activities.

C. Mathematical Modeling

The following notations are used to describe the DTCQTP:

- \( C_t \): Total cost of project (direct plus indirect).
- \( T_t \): Total duration of project, \( Qt \): Total quality of project, \( I_c \): Project indirect cost per time unit, \( Modes(i) \): Set of available execution modes for activity \( i \), \( C_{ik} \): Direct cost of activity \( i \) when performed the kth execution mode, \( tik \): Duration of activity \( i \) when performed the kth execution mode, \( qik \): Quality of activity \( i \) when performed the kth execution mode, \( y_{ik} \): Binary variable which is 1 when mode \( k \) is assigned to activity \( i \) and 0 otherwise, \( Defects Allowed \): Upper bound for Project quality. \( Tcpm \): Critical path duration obtained by critical bath method (CPM), if set of modes \( K = \{ k1, k2, ..., kn \} \) are assigned to activities;

Mixed integer programming is used for modeling DTCQTP:

\[
\text{Min } C_t = \sum_{i=0}^{N+1} \sum_{k \in Modes(i)} C_{ik} y_{ik} + I_c \cdot Tcpm \quad (12)
\]

\[
\text{Min } T_t = Tcpm \quad (13)
\]

Subject to:

\[
\sum_{i=0}^{N+1} \sum_{k \in Modes(i)} qik y_{ik} \leq \text{Defects Allowed} \quad (14)
\]

\[
\sum_{k \in Modes(i)} y_{ik} = 1. \quad (15)
\]

Objective functions (1) and (2) minimize the project's total costs and duration respectively. Constrain (3) enforces that the total quality of project does not bypass the desired level (upper bound). In (4) one and only one execution mode is assigned to each activity and equation (5) is sign constrains.

D. Problem Data Based Optimization (PDBO) Algorithm

PDBO algorithm is a single agent meta-heuristic algorithm that depends on possibility calculated from problem's data. PDBO assumes the problem is represented in the form of a graph \( G = (V, E) \), in which the set of nodes \( V \) represents the activities and modes, and the set of \( E \) represents edges that connects between activities and modes. For optimization problems, at each iteration, PDBO selects the first node \( n_i \) then depending on the best possibility values, it moves to the next adjacent node \( n_k \). After then, in order to increase the chance of selecting other nodes rather than node \( n_i \) in the next iteration, PDBO technique updates the possibility \( n_{i \rightarrow n_k} \) to be \( \text{Possibility}(n_{i \rightarrow n_k}) = \text{Possibility}(n_{i \rightarrow n_k}) + (\text{cost} / \alpha) \) where \( \alpha > 0 \).

Finally, after the last iteration solution found, in order to evaporate the Npossibilities, PDBO considers the parameter \( \beta \in [0, 1] \), such that \( \text{Possibility}(n_{i \rightarrow n_k}) = \text{Possibility}(n_{i \rightarrow n_k}) \rightarrow \beta \), where \( \beta \) is the evaporation rate (reduction rate) of possibility for virtual edge between \( n_i \) and \( n_k \).

PDBO manipulates the DTCQTP in the form of graph, in which the activities considered as nodes and the modes considered as virtual nodes connected to the activities through virtual edges fig 5. Possibility is calculated from problem's data (costs of activities). The PDBO algorithm for DTCQTP is shown in fig 7.

VI. EXAMPLE

Example with five task software programming project is considered fig 6, where activity 1 represents feasibility study and 2, 3, 4 and 5 represent requirements analysis, design, and code, documentation activities respectively and S and F are dummy activities. If bids offered from different subcontractors to perform specific activities in this project or this project will be developed in house, the modes take the forms as in table II below where task quality in each mode is estimated by COQUAMO and estimated software size is 25000 SLOC. In some large projects, many bids offered from many subcontractors to perform specific activities or tasks.

PDBO algorithm is implemented in c# and is tested and evaluated on CPU (Core i5) 3210 M, 2.50 GHz and 4GB RAM using Windows 7 as the operating system. Table III shows the results in terms of total quality \( Qt \), cost \( Ct \) and time \( Tt \) with Direct cost \( Cd \) of applying PDBO to this project using different quality bounds (Defects Allowed). This project has about 5^4 solutions. The parameters of PDBO and direct costs are included in the first row of table III. From table III and fig 8, the total cost and time are increased by minimizing quality bounds (minimizing defects) which means maximizing project quality. Fig 9 shows the processing time PT taken by PDBO algorithm to reach the solutions under different quality bounds.
TABLE II: EXECUTION MODES OF ACTIVITIES

<table>
<thead>
<tr>
<th>Activity number</th>
<th>Modes</th>
<th>Cost in $</th>
<th>Time in days</th>
<th>Estimated Defects</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>60</td>
<td>100</td>
<td>2.7</td>
</tr>
<tr>
<td>2</td>
<td>1</td>
<td>65</td>
<td>80</td>
<td>5.57940171679557</td>
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<tr>
<td>3</td>
<td>2</td>
<td>90</td>
<td>70</td>
<td>12.1229105965972</td>
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<tr>
<td>4</td>
<td>1</td>
<td>45</td>
<td>90</td>
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<tr>
<td>5</td>
<td>2</td>
<td>55</td>
<td>90</td>
<td>5.4</td>
</tr>
<tr>
<td>6</td>
<td>3</td>
<td>80</td>
<td>65</td>
<td>11.1588034335911</td>
</tr>
<tr>
<td>7</td>
<td>2</td>
<td>100</td>
<td>45</td>
<td>24.2458211931943</td>
</tr>
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</table>

TABLE III: FINAL OUTPUT OF PROGRAM

<table>
<thead>
<tr>
<th># of iterations=500, α=1000, β=0.02, indirect cost IC =20</th>
</tr>
</thead>
<tbody>
<tr>
<td>Qt</td>
</tr>
<tr>
<td>-----</td>
</tr>
<tr>
<td>51.987177338 4678</td>
</tr>
<tr>
<td>45.443668458 6663</td>
</tr>
<tr>
<td>39.684865025 0751</td>
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<tr>
<td>39.684865025 0751</td>
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<tr>
<td>31.66651459 8932</td>
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<tr>
<td>28.787159743 0976</td>
</tr>
</tbody>
</table>

Fig. 8. Time-cost-quality trade-off analysis

Fig. 9. Processing time (PT)

1. Set parameters $\alpha$, $\beta$ and read problem data
2. Calculate $Possibility(act, mods)$ from problem data , $i=1,...,Nc$, $k=1,...,Modes(act)$.
3. While (termination condition not met) do
4. For each activity $act_i$
5. Choose a mode $mod_k$ of activity $act_i$ in TCQTP with Minimum Possibilit.
6. Update possibility of virtual edge between $act_i$, and chosen mode $mod_k$ by $Possibility(act,mod_k) = Possibility (act,mod_k) + c(act,mod_k)/\alpha$.
7. End $i$
8. Find iteration solution
9. Evaporate a possibilities of virtual edges $e(act,mod_k)$ between all activities and their chosen modes at this iteration by $possibility(act,mod_k)=Possibility (act,mod_k) - \beta$.
10. Update a total solution
11. Evaporate possibilities of virtual edges $e(act,mod_k)$ that represent a total solution if no iteration solution exist by $possibility(act,mod_k)=Possibility (act,mod_k) - \beta$.
12. End while
13. Return a total solution

VII. CONCLUSION AND DISCUSSION

In this paper, task quality in each mode is estimated by COQUAMO and thus a total number of residual defects for all tasks in chosen modes reflect the quality of the software project quantitatively.

In DTCQTP, each project task/activity can be executed in one of several modes. The execution modes of any activity were assumed to be bids offered from different subcontractors to perform specific tasks or different approaches can be used to perform the tasks if the software project will be developed in house.

Solving the problem gave an optimal/nearly optimal solution in terms of time, cost, and quality of the project. By changing the allowable quality bound for the project and re-running the algorithm, other optimal /nearly optimal solutions could be obtained. Having these optimal/nearly optimal solutions and analyzing the environments needs, project managers could make decisions effectively.

To solve the problem, PDBO algorithm was introduced, which takes much less time to reach the optimal/nearly optimal solution under the allowable quality bound.
To achieve accurate estimation of task quality, COQUAMO need to be calibrated to origination database. In future work, we enhance COQUAMO accuracy by PDBO algorithm or another meta heuristic algorithm using NASA database.

REFERENCES


