Optimal Policy for a Deteriorating Inventory Model with Finite Replenishment Rate and with Price Dependant Demand Rate and Cycle Length Dependant Price

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Abstract— In this paper, an inventory model with finite and constant replenishment rate, price dependant demand rate, time value of money and inflation, finite time horizon, lead time and exponential deterioration rate and with the objective of maximizing the present worth of the total system profit is developed. Using a dynamic programming based solution algorithm, we are able to find the optimal sequence of the cycles and also to obtain different optimal selling prices, optimal order quantities and optimal maximum inventories for the cycles with unequal lengths, which have never been done before for this model. We also use a numerical example to show accuracy of the solution procedure.

Index Terms — Inventory control, Dynamic programming, Deteriorating items, Finite replenishment rate.

I. INTRODUCTION

Inventories generally can be classified into three meta-categories based on: 1-obsolecence 2-deterioration 3-no obsolescence/deterioration.

Obsolescence refers to inventories that become obsolete at a certain time, for instance because of rapid changes in technology or introduction of a new product. Deterioration refers to damage, spoilage, dryness, vaporization, etc. The products like fresh food (meat, fish, fruits, and vegetables), photographic films, batteries, human blood, photographic films, etc having a maximum usable lifetime are known as perishable products and the products like alcohol, gasoline, radioactive substances, lubricants, glues, paints, chemical ingredients, etc having no shelf-life at all are known as decaying products. No obsolescence/deterioration refers to inventories that their shelf life can be indefinite and hence they would fall under no obsolescence/deterioration category. In this paper we deal with the deteriorating inventories. There are many classifications for the deterioration. Ghare and Schrader [3] classify the deteriorating properties of inventory into three categories: 1-direct spoilage, e.g., vegetable, fruit and fresh food etc.; 2-physical depletion, e.g., gasoline and alcohol etc.; 3-deterioration such as radiation change, negative spoiling and loss of efficacy in inventory, e.g., electronic components and medicine. From another point of view, deterioration can also be classified by the timevalue or the items' life of inventory. Raafat [6] categorizes

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deterioration by the time-value of inventory:

1-utility constant: its utility does not change significantly as time passes within its valid usage period, e.g., liquid medicine; utility increasing: its utility increases as time passes

,e.g., some alcoholic drinks; 3- utility decreasing: its utility decreases as time passes, e.g., vegetables, fruits and fresh foods, etc. On the other hand, Nahmias [5] classifies deterioration by the items' life of inventory: 1-fixed lifetime: items' lifetime is pre-specified and its lifetime is independent of the deteriorating factors; therefore, it is called time independent deterioration. In fact, the utility of these items decreases during their lifetime, and when passing their lifetime, the item will perish completely and become of no value, e.g., milk, inventory in a blood bank, and food, etc. 2-random lifetime: there is no specified lifetime for these items. The lifetime for these items is assumed as a random variable, and its probability distribution could be gamma, weibull, exponential, etc. These Items are also called time dependent deteriorating items, e.g., electronic components, chemicals, and medicine, etc. The scope of this study covers those deteriorating items which are classified as utility decreasing (as regards as their time-value) and also as random lifetime (as regard as their lifetime).

Research in this area started with the work of Whitin [10] who considered fashion goods deteriorating at the end of prescribed storage period. An exponentially decaying inventory was developed first by Ghare and Schrader [3]. They observed that certain commodities shrink with time by a proportion which can be approximated by a negative exponential function of time.

Wee and Law [7] employed the concept of inflation and time value of money into the model where the demand is price dependent and shortage is allowed. The model considered a production environment with a finite replenishment rate. An optimization framework was presented to derive optimal production and pricing policies when the total net present value is maximized.

Wee and Wang [9] modeled a production-inventory system for deteriorating items with time varying demands and completely backlogged shortages. They constructed production schedules for this system using the traditional scheduling strategy in which each cycle starts with replenishment and ends with shortages.

Wee and Law [8] considered a deteriorating inventory model under time value of money and a deterministic inventory system with the price-dependent demand. They applied the discounted cash flows (DCF) approach for problem analysis. A heuristic approach was presented to

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derive the near optimal replenishment and pricing policy that tries to maximize the total net present value profit.

Zhou et al [11] extended the work of Wee and Wang [9] by permitting part of the backlogged shortages to turn into lost sales, which was assumed to be a function of currently backlogged amount. They also considered an alternative scheduling strategy in which each cycle of a schedule starts with a period of shortages and then followed by continuous replenishment. By examining the profit performance of the two scheduling strategies (i.e., "start with replenishment" versus "start with shortages"), they showed that their alternative strategy produces schedules with superior cost and profit values.

Chen [1], [2] used the dynamic programming method to find the optimal sequence of the cycles for his inventory models with the objective of minimizing the present worth of the total system cost and with time dependant demand rate, complete backlogging shortages and without considering the selling price and the revenue. We change and modify his procedure in order to solve our model.

Most of the researchers who worked in this area, assumed that selling price, order quantity, maximum inventory and cycle length are constant over the planning horizon (finite or infinite). In the present research, we develop a deteriorating inventory model with constant replenishment rate, price dependant demand rate, time value of money and inflation, finite time horizon, lead time and exponential deterioration rate with the objective of maximizing the present worth of the total system profit. Using a dynamic programming based solution algorithm, we are able to find the optimal sequence of the cycles and also to obtain different optimal selling prices, optimal order quantities and optimal maximum inventories for the cycles with unequal lengths, which have never been done before for this model. We also use a numerical example to show accuracy of the solution procedure.

II. MODEL DEVELOPMENT

The assumptions and the notations used in this model are: 1-The replenishment starts with zero inventory level at time *i*. The replenishment stops at time *j*, where the inventory attains its maximum level. During the period [i, j], the inventory level is affected by replenishment, demand and deterioration of items. Due to reasons of demand and deterioration of items, the inventory level gradually diminishes during the period [j, k] and ultimately falls to zero at time *k*.

2-The demand rate, D(s) is a linear, non-negative, continuous, convex, decreasing function of the selling price, *s*. D(s) = a - bs, where *a* and *b* are both positive constants. 3-A single item is considered, whose time to deterioration is distributed as exponential. $f(t) = \lambda e^{-\lambda t}$, where λ is a positive constant (deterioration rate) and t > 0.

4-The inflation rate (denoted by *m*) and the discount rate (denoted by *r*) representing the time value of money are constant. Therefore, the net discount rate of inflation is constant: R = r - m. The present worth of X_t , is: $PW(X_t) = X_t e^{-Rt}$, $t \ge 0$. where X_t is the value of *X* at time *t*. Notice that both the internal (company) inflation rate and the external (general economy) inflation rate can be estimated by a unique inflation rate.

5-The replenishment rate is finite and constant, P, where P > D(s).

6-Lead time (*Lt*) is considered but is limited to $0 \le Lt \le (k-i)$, because of the inequality of the cycle lengths.

7-Holding cost applies to good units only.

8-Deterioration of units occurs only when the item is effectively in stock, and there is no repair or replacement of deteriorated units during planning horizon.

9-Shortages are not allowed.

In addition, the following notations are used in this paper: $I_i(t)$: inventory level at time *t*, where $i \le t \le j$.

 $I_{2}(t)$: inventory level at time t, where $j \le t \le k$.

Q: ordering quantity per cycle.

A: ordering cost per cycle at time zero.

C: unit purchasing cost at time zero, where C < S.

h: unit holding cost per unit of time at time zero.

H: planning horizon.

The inventory level can be described by the following differential equations:

$$\frac{\delta I_1(t)}{\delta t} + \lambda I_1(t) = P - D(s) \quad i \le t \le j \quad (1)$$
$$\frac{\delta I_2(t)}{\delta t} + \lambda I_2(t) = -D(s) \quad j \le t \le k \quad (2)$$

By applying the boundary conditions, multiplying e^{i} on the both sides of (1), (2) and integrating by part, we have:

$$I_{1}(t) = (P - D(s))e^{-\lambda(t-i)} \int_{i}^{t} e^{\lambda(u-i)} du =$$

$$(P - D(s))e^{-\lambda t} \int_{i}^{t} e^{\lambda u} du \qquad i \le t \le j$$

$$I_{2}(t) = D(s)e^{-\lambda(t-i)} \int_{t}^{k} e^{\lambda(u-i)} du = D(s)e^{-\lambda t} \int_{t}^{k} e^{\lambda u} du \qquad j \le t \le k$$

And because the maximum inventory takes place at time j, we have:

$$I_{\max} = (P - D(s))e^{-\lambda j} \int_{i}^{j} e^{\lambda u} du = D(s)e^{-\lambda j} \int_{j}^{k} e^{\lambda u} du$$
(3)

After simplification, we have:

$$P.e^{-\lambda j}\int_{i}^{j}e^{\lambda u}du=D(s)e^{-\lambda j}\int_{i}^{k}e^{\lambda u}du$$

By solving the above equation with respect to *j*, we have:

$$j = \frac{Ln[\lambda \frac{D(s)}{P} \int_{i}^{k} e^{\lambda u} du + e^{\lambda i}]}{\lambda} = \frac{Ln[\frac{D(s)}{P} (e^{\lambda k} - e^{\lambda i}) + e^{\lambda i}]}{\lambda} \quad (4)$$

Ordering quantity for each cycle can be obtained as follows: all of the deteriorated items during the cycle + all of the demands during the cycle.

$$Q = D(s) \int_{i}^{k} du + \int_{i}^{j} \lambda I_{1}(t) dt + \int_{j}^{k} \lambda I_{2}(t) dt =$$
(5)

$$D(s)(k-i) + (P-D(s))\mathcal{A}\int_{i}^{j}\int_{i}^{t} e^{\lambda(u-t)} du dt + D(s)\mathcal{A}\int_{j}^{k}\int_{t}^{k} e^{\lambda(u-t)} du dt$$

Reorder point in each cycle is (obviously, the order quantity would be the next cycle's order quantity):

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If
$$0 \le Lt \le (k-j)$$
 then $RoP = \int_{k-Lt}^{k} D(s)e^{\lambda(u-(k-Lt))} du$ (6)

If
$$(k-i) \ge Lt \ge (k-j)$$
 then

$$RoP = \int_{i}^{k-L} (P-D(s))e^{\lambda(u-(k-L))} du$$
(7)

Notice that if Lt = (k - j) then either one is usable.

The present worth of the relevant costs and the revenue for each cycle are as follows:

The Present worth of the holding cost:

$$Hc = h \int_{i}^{j} e^{-Rt} I_{1}(t) dt + h \int_{j}^{k} e^{-Rt} I_{2}(t) dt =$$

$$h.(P - D(s)) \int_{i}^{j} e^{-Rt} e^{-\lambda t} \int_{i}^{t} e^{\lambda u} du dt + h.D(s) \int_{j}^{k} e^{-Rt} e^{-\lambda t} \int_{t}^{k} e^{\lambda u} du dt$$

The present worth of the purchasing cost: $Pc = \int_{i}^{i} P.c.e^{-Ru} du$

The present worth of the ordering cost: $Oc = A.e^{-R(i-h)}$

The present worth of the revenue: $R = s \int_{a}^{b} D(s) e^{-Ru} du$

Now we can compute the present worth of the total system profit as follows: TP = P - HC - PC - CC

$$IP = R - HC - PC - OC$$

$$TP = s.D(s) \int_{i}^{k} e^{-Ru} du - h.(P - D(s)) \int_{i}^{j} e^{-Rt} e^{-\lambda t} \int_{i}^{t} e^{\lambda u} du dt - h.(P - D(s)) \int_{i}^{j} e^{-Rt} e^{-\lambda t} \int_{i}^{t} e^{\lambda u} du dt - h.(P - D(s)) \int_{i}^{j} e^{-Ru} du - A.e^{-R(i - lt)}$$
(8)

In order to determine the optimal selling price for each cycle we differentiate (8) with respect to s and we set the result equal to zero. Then, we have:

$$\frac{\partial TP}{\partial s} = (a - 2b.s) \int_{i}^{k} e^{-Ru} du + b[-h] \int_{i}^{j} e^{-Rt} e^{-\lambda t} \int_{i}^{l} e^{\lambda u} du dt + h \int_{j}^{k} e^{-Rt} e^{-\lambda t} \int_{i}^{k} e^{\lambda u} du dt] = 0$$

$$\Rightarrow s = (a \int_{i}^{k} e^{-Ru} du + b[-h] \int_{i}^{j} e^{-Rt} e^{-\lambda t} \int_{i}^{t} e^{\lambda u} du dt + h \int_{j}^{k} e^{-Rt} e^{-\lambda t} \int_{i}^{t} e^{\lambda u} du dt + h \int_{j}^{k} e^{-Rt} e^{-\lambda t} \int_{i}^{k} e^{\lambda u} du dt] / 2b \int_{i}^{k} e^{-Ru} du$$
(9)

To satisfy the optimally condition we must have $\frac{\partial^2 TP}{\partial^2 s} < 0.$

$$\frac{\partial^2 TP}{\partial^2 s} = -2b \int_{i}^{j} e^{-Ru} du < 0.$$

Then, the optimally condition is strongly satisfied, and our obtained s will maximize the present worth of the total system profit and *TP* is also concave with respect to s.

III. SOLUTION PROCEDURE

For obtaining the optimal sequence of the cycles, as we mentioned in the introduction section, we change and modify the dynamic programming model and algorithm used by Chen [1], [2], as follows:

$$TP_k = \underset{0 \le i < j < k \le H}{Max} \{TP_i + W(i, s^*, j^*, k)\}$$
$$TP_0 = 0$$

This recursive procedure works in a forward fashion to determine the maximal present worth of the total profit over the time horizon, H. Where, $0 \le i < j < k \le H$, TP_k is the optimal present worth of the profit until time *k* so that TP_H is the optimal present worth of the total system profit and $W(i, s^*, j^*, k)$ is the optimal present worth of the profit in cycle [i, k] knowing s^* which is the optimal *s* in cycle [i, k] and *j* which is the point where the replenishment ends in cycle [i, k].

The dynamic programming based solution algorithm to solve this model is:

Step 1: Enter parameters: $P, c, H, h, \lambda, a, b, R, lt, A$.

Step 2: Let $i = 0, TP_0 = 0$ For K = 1: H Obtain *j*, *S* from (4), (9); (solve together) Obtain $W(i, s^*, j^*, k)$ from (8) Let $TP_{k} = W(i, s^{*}, j^{*}, k), s^{*}_{i,k} = s, j^{*}_{i,k} = j, T_{k} = i$ End Step 3: For k = 2: H For i = 1: k-1Obtain j, S from (4), (9); (solve together) Obtain $W(i, s^*, j^*, k)$ from (8) If $TP_{i} < TP_{i} + W(i, s^{*}, j^{*}, k)$ Then $TP_{k} = TP_{i} + W(i, s^{*}, j^{*}, k), s^{*}_{i,k} = s, j^{*}_{i,k} = j, T_{k} = i$ End End Step 4: Let k = HWhile k > 0Cycle time = $[T_k, k]$ Optimal price = $s^*_{T_k,k}$ Replenishment ending time= $j^*_{T_k,k}$ Optimal maximum inventory= $I_{\max T_{i},k}^{*}$, (from (3)) Optimal order quantity = $Q^*(T_i, s^*, j^*, k)$, (from (5)) Optimal reorder point: $RoP_{T_{k,k}}$, (if $0 \le Lt \le (k-j)$ from (6), if $(k - i) \ge Lt > (k - j)$ from (7) and if Lt > (k - i)) then 'Lead time not acceptable') Cumulated total profit = TP_{μ} Let $k = T_k$ End

End

Where, T_k is the starting point of the last replenishment cycle from time zero to time k (i.e., $[T_k, k]$) $T_k = 0,1,2,...,k-1$ and k = 1,2,...,H, $s^*_{T_k,k}$ is the optimal selling price in cycle $[T_k, k]$, $j^*_{T_k,k}$ is the time in cycle $[T_k, k]$ where the replenishment ends, $Q^*(T_j, s^*, j^*, k)$ is the optimal order quantity in cycle $[T_k, k]$ knowing $s^*_{T_k,k}$ and $RoP_{T_k,k}$ is the optimal reorder point in cycle $[T_k, k]$.

IV. NUMERICAL EXAMPLE

In this section, we use a numerical example to show accuracy of the model and the solution procedure. The given parameters' values are: Proceedings of the World Congress on Engineering and Computer Science 2008 WCECS 2008, October 22 - 24, 2008, San Francisco, USA

$$a = 50, b = 0.9, R = 0.1, \lambda = 0.05, h = 0.5, c = 8, A = 100,$$

H = 10, P = 95, Lt = 0.5.

We solve the model by MATLAB 7.1 to obtain the necessary results and they can be seen in Table I. Then, we run the proposed algorithm to obtain the optimal sequence of the cycles. After running the algorithm by MATLAB 7.1, we have two cycles with lengths of 2 and two cycles with lengths of 3, with the present worth of the total system profit equal to 2565.4, as follows:

 $TP_{10} = W(8, s^*, j^*, 10) + W(6, s^*, j^*, 8) + W(3, s^*, j^*, 6) +$

 $W(0, s^*, j^*, 3) = 2565.4$

The optimal price and the optimal order quantity for the optimal cycles are (for optimal maximum inventory, optimal reorder point, optimal replenishment ending time and more detail you can refer to Table I):

$$s^*_{6,8} = s^*_{8,10} = 27.8936, s^*_{0,3} = s^*_{3,6} = 27.9491$$

 $Q^*_{6,8} = Q^*_{8,10} = 51.6576, Q^*_{0,3} = Q^*_{3,6} = 78.7628$

Table I
Software output for the numerical example

i	k	$j^*_{T_k,k}$	S [*] <i>i,k</i>	$W(i, s^*, j^*, k)$	$I^*{}_{\max}$	Q^*	RoP^*
0	1	0.2675	27.8364	350.6848	18.6131	25.4108	12.6309
0	2	0.5438	27.8936	734.2736	37.6066	51.6576	12.6048
1	2	1.2675	27.8364	317.3127	18.6131	25,4108	12.6309
0	3	0.8291	27.9491	1.0527e+003	56.9745	78.7628	12.5795
1	3	1.5438	27.8936	664.3982	37.6066	51.6576	12.6048
2	3	2.2675	27.8364	287.1164	18.6131	25.4108	12.6309
0	4	1.1237	28.0024	1.3123e+003	76.7101	106.7502	12.5552
1	4	1.8291	27.9491	952.5110	56.9745	78.7628	12.5795
2	4	2.5438	27.8936	601.1724	37.6066	51.6576	12.6048
3	4	3.2675	27.8364	259.7937	18.6131	25.4108	12.6309
0	5	1.4278	28.0534	1.5191e+003	96.8060	135.6444	12.5320
1	5	2.1237	28.0024	1.1875e+003	76.7101	106.7502	12.5552
2	5	2.8291	27.9491	861.8676	56.9745	78.7628	12.5795
3	5	3.5438	27.8936	543.9632	37.6066	51.6576	12.6048
4	5	4.2675	27.8364	235.0710	18.6131	25.4108	12.6309
0	6	1.7418	28.1015	1.6782e+003	117.2538	165.4713	12.5100
1	6	2.4278	28.0534	1.3745e+003	96.8060	135.6444	12.5320
2	6	3.1237	28.0024	1.0745e+003	76.7101	106.7502	12.5552
3	6	3.8291	27.9491	779.8500	56.9745	78.7628	12.5795
4	6	4.5438	27.8936	492.1983	37.6066	51.6576	12.6048
5	6	5.2675	27.8364	212.7011	18.6131	25.4108	12.6309
0	7	2.0659	28.1465	1.7946e+003	138.0444	196.2578	12.4895
1	7	2.7418	28.1015	1.5185e+003	117.2538	165.4713	12.5100
2	7	3.4278	28.0534	1.2437e+003	96.8060	135.6444	12.5320
3	7	4.1237	28.0024	972.2056	76.7101	106.7502	12.5552
4	7	4.8291	27.9491	705.6375	56.9745	78.7628	12.5795
5	7	5.5438	27.8936	445.3594	37.6066	51.6576	12.6048
6	7	6.2675	27.8364	192.4599	18.6131	25.4108	12.6309
0	8	2.4003	28.1881	1.8727e+003	159.1675	228.0320	12.4706
1	8	3.0659	28.1465	1.6238e+003	138.0444	196.2578	12.4895
2	8	3.7418	28.1015	1.3740e+003	117.2538	165.4713	12.5100
3	8	4.4278	28.0534	1.1254e+003	96.8060	135.6444	12.5320
4	8	5.1237	28.0024	879.6880	76.7101	106.7502	12.5552
5	8	5.8291	27.9491	638.4872	56.9745	78.7628	12.5795
6	8	6.5438	27.8936	402.9779	37.6066	51.6576	12.6048
7	8	7.2675	27.8364	174.1449	18.6131	25.4108	12.6309
0	9	2.7455	28.2258	1.9165e+003	180.6120	260.8229	12.4534
1	9	3.4003	28.1881	1.6945e+003	159.1675	228.0320	12.4706
2	9	4.0659	28.1465	1.4693e+003	138.0444	196.2578	12.4895
3	9	4.7418	28.1015	1.2432e+003	117.2538	165.4713	12.5100
4	9	5.4278	28.0534	1.0183e+003	96.8060	135.6444	12.5320
5	9	6.1237	28.0024	795.9746	76.7101	106.7502	12.5552
6	9	6.8291	27.9491	577.7271	56.9745	/8./628	12.5795
1	9	7.5438	27.8936	364.6295	37.6066	51.6576	12.6048
8	9	8.2675	27.8364	157.5728	18.6131	25.4108	12.6309
0	10	3.1017	28.2593	1.929/e+003	202.3652	294.6604	
1	10	3.7455	28.2258	1.7341e+003	180.6120	260.8229	
2	10	4.4003	28.1881	1.5352e+003	139.16/5	228.0320	
5	10	5.0659	28.1465	1.3295e+003	138.0444	196.2578	
4	10	5./418	28.1015	1.1249e+003	117.2538	105.4/13	
3	10	0.4278	28.0534	921.3042	90.8000	105.0444	
0	10	7.9201	28.0024	720.2276	/0./101	106./502	
/	10	/.8291	27.9491	322.7491	27.6066	/8./628	
ð	10	8.5458 0.2675	27.8930	529.9304	37.0000	31.03/0	
9	10	1 9/0/5	1 2/ 8364	142.5778	1 18.6131	274108	

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V. CONCLUSIONS

A single item inventory model with constant replenishment rate, exponential deteriorating rate, price dependant demand rate, time value of money and inflation, lead time, finite time horizon and with the objective of maximizing the present worth of the total system profit, was developed in this paper. We used a numerical example to solve the model and we concluded different optimal selling prices, optimal order quantities, optimal maximum inventories and optimal time fractions of replenishment for the cycles with unequal lengths. Then, we ran a dynamic programming based algorithm to obtain the optimal sequence of the cycles and the results showed that the optimal system has two cycles with lengths of 2 and two cycles with lengths of 3.

In addition, according to the numerical analysis, we can conclude that with increasing/decreasing of the cycle length; present worth of cycle profit, optimal selling price, optimal order quantity, optimal maximum inventory and optimal time fraction of replenishment will increase/decrease. Even with reasonable changes in the parameters' values, we will have the above results satisfied (we have tested this by changing the parameters' values).

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